# Distributed Binary Detection Over Fading Channels: Cooperative and Parallel Architectures

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Abstract—This paper considers the problem of binary distributed detection of a known signal in correlated Gaussian sensing noise in a wireless sensor network, where sensors are restricted to using the likelihood ratio test (LRT) and communicating with the fusion center (FC) over bandwidth-constrained channels that are subject to fading and noise. To mitigate the deteriorating effect of fading encountered in the conventional parallel fusion architecture, in which sensors directly communicate with the FC, we propose new fusion architectures that enhance the detection performance, via harvesting cooperative gain (so-called "decision diversity gain"). In particular, we propose 1) cooperative fusion architecture with Alamouti's space-time coding scheme at sensors, 2) cooperative fusion architecture with signal fusion at sensors, and 3) parallel fusion architecture with local threshold changing at sensors. For these schemes, we derive the LRT and majority fusion rules at the FC and provide upper bounds on the average error probabilities for homogeneous sensors, subject to uncorrelated Gaussian sensing noise, in terms of signal-to-noise ratio (SNR) of communication and sensing channels. Our simulation results indicate that when the FC employs the LRT rule, except for low communication SNR and moderate/high sensing SNR, performance improvement is feasible with the new fusion architectures. When the FC utilizes the majority rule, such an improvement is possible, except for high sensing SNR.

*Index Terms*—Correlation, distributed detection, diversity, error floor, fusion and sensor rule, parallel architecture, space-time coding.

### I. INTRODUCTION

T HE problem of distributed detection with the fusion center (FC; so-called classical parallel fusion architecture) has a long and rich history, where each local detector (sensor) processes its observation locally and independently and passes its local binary decision to the FC. The main assumption in the classical works is that the bandwidth-constrained communication channels are *error free*, and thus, the reliability of the final

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decision at the FC is determined by the reliability of the local binary decisions. However, wireless channels are inherently error prone, due to noise and fading. An integrated approach of distributed detection over noisy fading channels was considered in [1]-[12], in which the sensors send their modulated local binary decisions to the FC, and the FC employs a fusion rule, incorporating channel state information (CSI), to improve the reliability of the final decision at the FC. The performance of this integrated distributed detection is ultimately limited by the communication bounds. A common thread in the schemes discussed in [1]–[12] is that they are noncooperative, i.e., there is no information exchange among the sensors. Cooperative wireless communication [13], [18] has been proven to significantly enhance performance in the presence of fading, via invoking spatial diversity, that leads to the mitigation of the detrimental fading effects [13], [18]. Motivated by the promises of cooperative communication, we propose a new class of integrated distributed detection, which harvests cooperative gain (enabled by, at most, 1-bit information exchange among onehop neighboring nodes) and improves the performance of the integrated distributed detection [1], [2] in the presence of fading, via allowing each sensor to send (at most) two information bits to the FC and assuming identical transmit power per node. In particular, we propose three schemes: 1) cooperative fusion architecture with Alamouti's space-time coding (STC) scheme at sensors, in which neighboring sensors exchange one information bit, and each sensor sends two information bits to the FC; 2) cooperative fusion architecture with signal fusion at sensors, in which neighboring sensors exchange one information bit, and each sensor sends one information bit to the FC; and 3) parallel fusion architecture with local threshold changing at sensors, in which neighboring sensors do not exchange information, and each sensor sends two information bits to the FC. To describe the proposed schemes, suppose  $S_1$  and  $S_2$  are two designated cooperative partners. In scheme 1), rather than transmitting their local decisions directly to the FC,  $S_1$  and  $S_2$  are coordinated to form a transmit cluster, such that they first exchange their local decisions and apply Alamouti's scheme [19], [20] for transmitting the decisions to the FC. Different from most literature on distributed STC, which assume that a node acts as a relay only for error-free reception [21], we consider the fact that the channels between  $S_1$  and  $S_2$  are subject to errors, due to noise and fading. In scheme 2), similar to scheme 1),  $S_1$  and  $S_2$  exchange their local decisions. Instead of applying Alamouti's scheme, however, each node updates its decision, via optimally fusing its observation with the received signal from its cooperative partner. Updated decisions are transmitted

to the FC. In scheme 3), different from schemes 1) and 2), there is no explicit information exchange between  $S_1$  and  $S_2$ . Each node forms two decisions, where one decision is made based on its observation only, and the other decision is obtained based on optimally fusing its observation with its guess of the decision of its cooperative partner.  $S_1$  and  $S_2$  apply Alamouti's scheme for transmitting these decisions to the FC. For these three schemes, we provide the likelihood ratio test (LRT) and majority fusion rules at the FC. The average (over fading) error probability of the decision error at the FC for these schemes depends on the signal-to-noise ratio (SNR) of communication (channels between cooperative partners as well as between the nodes and the FC) and sensing channels (channels between the target and the nodes). For the proposed schemes, we derive upper bounds on the average error probability and investigate how a node should allocate its transmit power for communicating with its cooperative partner and the FC, such that the error is minimized. These results enable us to quantity the cooperative gain offered by the proposed schemes, with respect to the schemes in [1] and [2], assuming identical transmit power per node.

Information exchange among the nodes for consensus-based distributed detection without the FC has been previously studied (examples are [14]–[17]). In [14]–[16], a consensus-based distributed detection system is considered, where the sensors successively update and broadcast their local binary decisions over error free links [14]–[16]. In [17], each sensor successively updates its continuous-valued decision variable and passes it to its neighbors over random links without a bandwidth constraint. Distributed detection in networks with feedback has been previously studied (examples are [23]-[25]) from the information-theoretic perspective, focusing on the asymptotic regime (i.e., networks with a large number of sensors) and quantifying performance in terms of error exponents. These works consider a variety of feedback architectures, including two-message feedback architectures, where each sensor sends its first message to the FC, based on its own observation, and sends its second message, based on the additional information provided by the FC through feedback, where the feedback contains (functions of) the messages generated by (some) other sensors. Different from these works, we consider networks with a finite number of sensors, without feedback from the FC, where (at most) 1-bit information exchange is allowed between two cooperative sensors. Moreover, we relax the errorfree communication constraint, by considering fading effects during the information exchange phase, i.e., we assume that a sensor knows the decision of its partner with limited reliability, that is dictated by the quality of the intersensor communication channel. Perhaps the most related work is that in [22], in which each sensor communicates its local binary decision to its neighbors, and the sensors communicate their updated binary decisions to the FC. The local decision rules and the fusion rule at the FC are all majority rules, and communication channels are assumed to be error free in [22]. To the best of our knowledge, for parallel fusion architecture, no prior work has studied the impact of local (limited) information exchange on enhancing the performance of the integrated distributed detection systems [1]-[12] operating in a noisy fading environment. This paper is organized as follows. Section II

introduces our sensing model and overviews the integrated distributed detection schemes in [1] and [2] for this work to be self-explanatory. Sections III–V describe schemes 1)–3), respectively, and provide local decision and fusion rules at the sensors and the FC. Section VI provides the performance analysis. Our numerical results are presented in Section VII. Concluding remarks are in Section VIII.

### **II. BASIC MODELS**

### A. Sensing Model

We consider the binary hypothesis testing problem of detecting a known signal in correlated Gaussian noise based on measurements  $x_k$  at K distributed sensors. The *a priori* probabilities of two hypotheses  $\mathcal{H}_0, \mathcal{H}_1$  are denoted by  $\pi_0, \pi_1$ , respectively. The FC is tasked with determining whether the unknown hypothesis is  $\mathcal{H}_0$  or  $\mathcal{H}_1$ , based on the information collected from the K sensors. The measurement  $x_k$  of sensor  $\mathcal{S}_k$  under  $\mathcal{H}_0$  and  $\mathcal{H}_1$ , respectively, are  $x_k = w_k$  and  $x_k =$  $1 + w_k$  for  $k = 1, \ldots, K$ , where sensing noise  $w_k$  is zero mean with variance  $\sigma_{w_k}^2$ . The spatial correlation between noise terms  $w_i, w_j$  is characterized with the correlation coefficient  $\rho_{ij}$ . We assume that the sensors are grouped into S = K/2 distinct pairs of cooperative partners  $(S_i, S_j)$ , where  $S_i$  knows  $\sigma_{w_i}^2$ ,  $\sigma_{w_i}^2$ , and  $\rho_{ij}$  only and is restricted to use the LRT. The FC employs the LRT, when all sensing noise variances and pairwise correlation coefficients are available at the FC. In the absence of this knowledge, the FC uses the majority rule.

#### B. Classical Parallel Fusion Architecture

Each sensor  $S_k$  makes a local binary decision  $u_k \in \{1, -1\}$ based on its measurement  $x_k$ . The local decisions +1 and -1 correspond to the hypotheses  $\mathcal{H}_1$  and  $\mathcal{H}_0$ , respectively. The local detection performance of  $S_k$  is characterized with  $P_{d_k} =$  $P(u_k = 1 | \mathcal{H}_1)$  and  $P_{f_k} = P(u_k = 1 | \mathcal{H}_0)$ . To form  $u_k$ , sensor  $\mathcal{S}_k$  applies the LRT  $(f(x_k|\mathcal{H}_1)/f(x_k|\mathcal{H}_0)) \geq_{u_k=-1}^{u_k=1} (\pi_0/\pi_1).$ For the sensing model in Section II-A, the LRT is reduced to  $x_k \gtrsim_{u_k=-1}^{u_k=1} \tau_k$ , where  $\tau_k = 0.5 + \sigma_{w_k}^2 \ln(\pi_0/\pi_1)$ . Moreover,  $P_{d_k} = Q((\tau_k - 1)/\sigma_{w_k})$ , and  $P_{f_k} = Q(\tau_k/\sigma_{w_k})$ , where Q(x)denotes the Q-function.<sup>1</sup> The local decisions  $u_k$  are transmitted over orthogonal channels subject to noise and fading to the FC. When  $S_k$  sends  $u_k$ , the received signal at the FC is  $y_k =$  $u_k h_k + v_k$ , where  $h_k \sim \mathcal{CN}(0, \sigma_{h_k}^2), v_k \sim \mathcal{CN}(0, \sigma_v^2)$ , with  $h_k$ representing the fading channel coefficient corresponding to the channel between  $S_k$  and the FC. The channel variance is  $\sigma_{h_k}^2 = \mathcal{PG}/d_k^{\varepsilon}$ , where  $\mathcal{P}$  represents the transmit power of  $\mathcal{S}_k, d_k$  denotes the distance between  $\mathcal{S}_k$  and the FC,  $\varepsilon$  is the path-loss exponent, and  $\mathcal{G}$  is a constant that depends on the antenna gains and the wavelength. We assume that  $\varepsilon$  and  $\mathcal{G}$  are identical for all links. The term  $v_k$  is the receiver noise at the FC. We assume that  $v_k$  and  $h_k$  are independent.

<sup>&</sup>lt;sup>1</sup>Considering the local LRT at sensor k, the distributions  $f(x_k | \mathcal{H}_\ell)$  for  $\ell = 0, 1$  and  $P_{d_k}, P_{f_k}$  values depend on the distribution of sensing noise. Thus, the simplified form of the local rule and  $P_{d_k}, P_{f_k}$  expressions would change for non-Gaussian  $w_k$ . Given the joint probability density function (pdf)  $f(w_1, \ldots, w_K)$  at the FC, the expressions  $\Lambda$  in (1) for dependent and (2) for independent  $w_k$ 's remain unchanged.

We define  $\bar{\gamma}_{h_k}^2 = \sigma_{h_k}^2/\sigma_v^2$  as the average received SNR corresponding to node  $S_k$ -FC communication channel. Relying on the received signals  $y_k$  and the availability of CSI  $h_k$ , the FC forms the LRT  $\Lambda = f(y_1, \ldots, y_K | \mathcal{H}_1)/f(y_1, \ldots, y_K | \mathcal{H}_0)$ , where  $f(y_1, \ldots, y_K | \mathcal{H}_\ell)$  indicates the joint pdf of  $y_1, \ldots, y_K$  given the hypothesis  $\mathcal{H}_\ell$ , to make the final decision. In particular, the FC decides on  $\mathcal{H}_1$  when  $\Lambda > \pi_0/\pi_1$  and decides on  $\mathcal{H}_0$  otherwise. Considering that  $\mathcal{H}_\ell \to u_1, \ldots, u_K \to y_1, \ldots, y_K$  form a Markov chain and the fact  $f(y_1, \ldots, y_K | u_1, \ldots, u_K) = \prod_{k=1}^K f(y_k | u_k)$ , we can simplify  $\Lambda$  as follows:

$$\Lambda = \frac{\sum_{u_1} \cdots \sum_{u_K} \left( \prod_{k=1}^K f(y_k | u_k) \right) P(u_1, \dots, u_K | \mathcal{H}_1)}{\sum_{u_1} \cdots \sum_{u_K} \left( \prod_{k=1}^K f(y_k | u_k) \right) P(u_1, \dots, u_K | \mathcal{H}_0)}.$$
(1)

Focusing on the terms  $f(y_k|u_k)$  in (1), we note that given  $h_k, u_k$ , we have  $y_k \sim C\mathcal{N}(u_k h_k, \sigma_v^2)$ . Considering the term  $P(u_1, \ldots, u_K | \mathcal{H}_\ell)$  in (1), we note that it depends on the characteristics of the sensing channel noise described in Section II-A. When Gaussian sensing noise terms are uncorrelated  $\rho_{ij} = 0$  for all  $i \neq j$ , this term is simplified to  $\prod_{k=1}^K P(u_k | \mathcal{H}_\ell)$ , and (1) is reduced to [1], [2]

$$\Lambda = \prod_{k=1}^{K} \frac{P_{d_k} f(y_k | u_k = 1) + (1 - P_{d_k}) f(y_k | u_k = -1)}{P_{f_k} f(y_k | u_k = 1) + (1 - P_{f_k}) f(y_k | u_k = -1)}.$$
 (2)

When the parameters of sensing channels are unavailable at the FC, the FC cannot apply the optimal LRT in (1) or (2). Alternatively, the FC demodulates the channel inputs  $u_k$  using  $y_k$  for k = 1, ..., K and applies the majority rule to the demodulated symbols to reach the final decision, i.e., if the sum of the demodulated symbols is positive, the FC decides on  $\mathcal{H}_1$  and, otherwise, decides on  $\mathcal{H}_0$ .

### III. COOPERATIVE FUSION ARCHITECTURE WITH SPACE–TIME CODING AT SENSORS

# *A. Internode Communication Channel Model and Local Decision Rules at Sensors*

Suppose nodes  $S_i$  and  $S_j$  are within a pair, that is,  $S_i$  and  $S_j$  are cooperative partners. Each sensor makes a decision based on its measurement. The nodes exchange their decisions over orthogonal channels subject to noise and fading. Let  $u_i$  denote the decision made at  $S_i$  based on  $x_i$ . Sensor  $S_i$  transmits  $\sqrt{1 - \alpha u_i}$ , where  $0 < \alpha < 1$  is a power normalization factor, to assure that the total transmit power of nodes in this architecture remains the same as that of the classical parallel fusion architecture in Section II-B. When  $S_i$  transmits  $\sqrt{1 - \alpha u_i}$ , the received signal at  $S_j$  is

$$r_{ij} = \sqrt{1 - \alpha} u_i g_{ij} + \eta_{ij}$$
  
where  $g_{ij} \sim C\mathcal{N}\left(0, \sigma_{hs_{ij}}^2\right), \quad \eta_{ij} \sim C\mathcal{N}\left(0, \sigma_{\eta}^2\right).$  (3)

 $g_{ij}$  represents the fading channel coefficient from  $S_i$  to  $S_j$ , and  $\eta_{ij}$  is the receiver noise at  $S_j$ . The channel variance is  $\sigma_{h_{S_{ij}}}^2 = \mathcal{PG}/d_{ij}^{\varepsilon}$ , where  $d_{ij}$  denotes the distance between  $S_i$  and  $S_j$ . Noise terms  $\eta_{ij}$ ,  $\eta_{ji}$  and channel coefficients  $g_{ij}$ ,  $g_{ji}$  are independent, and noise terms are independent and identically distributed (i.i.d.) across all pairs. Upon receiving  $r_{ij}$ ,  $S_j$  demodulates the channel input  $u_i$ , using the knowledge of  $g_{ij}$ , i.e.,

$$\hat{u}_i = \operatorname{sgn}\left(\mathfrak{Re}\left(\frac{r_{ij}}{g_{ij}}\right)\right).$$
 (4)

The pair  $(S_i, S_j)$  sends the information  $u_i, \hat{u}_i, u_j, \hat{u}_j$  to the FC in two consecutive time slots, exploiting Alamouti's STC scheme. In particular, in the *n*th slot,  $S_i$  and  $S_j$  simultaneously send  $\sqrt{(\alpha/2)}u_i$  and  $\sqrt{(\alpha/2)}u_j$ , respectively. In the (n + 1)th time slot,  $S_i$  and  $S_j$  simultaneously send  $-\sqrt{(\alpha/2)}\hat{u}_j$  and  $\sqrt{(\alpha/2)}\hat{u}_i$ , respectively. Considering the definitions of channel variances, we note that, effectively,  $S_i$  spends  $(1 - \alpha)\mathcal{P}$  and  $\alpha\mathcal{P} = (\alpha/2)\mathcal{P} + (\alpha/2)\mathcal{P}$ , respectively, for internode and sensor–FC communication.

# *B.* Node–FC Communication Channel Model and Fusion Rule at the FC

Let  $y_{ij}(n)$  and  $y_{ij}(n + 1)$  denote the received signals at the FC corresponding to the pair  $(S_i, S_j)$  during two consecutive time slots. We have

$$y_{ij}(n) = \sqrt{\frac{\alpha}{2}} (u_i h_i + u_j h_j) + v_{ij}(n)$$
$$y_{ij}(n+1) = \sqrt{\frac{\alpha}{2}} (-\hat{u}_j h_i + \hat{u}_i h_j) + v_{ij}(n+1)$$
(5)

where  $h_{i(j)} \sim C\mathcal{N}(0, \sigma_{h_{i(j)}}^2), v_{ij}(n), v_{ij}(n+1) \sim C\mathcal{N}(0, \sigma_v^2)$ . The term  $v_{ij}(n)$  is the receiver noise at the FC during the *n*th time slot. We assume that noise terms  $v_{ij}(n), v_{ij}(n+1)$  and channel coefficients  $h_i, h_j$  are independent, and noise terms are i.i.d. across all pairs. Taking similar processing as Alamouti decoding [20], the FC first forms  $z_i, z_j$  using  $y_{ij}(n), y_{ij}^*(n+1)$  as follows:

$$\begin{split} \begin{bmatrix} z_i \\ z_j \end{bmatrix} &= \begin{bmatrix} h_i^* & h_j \\ h_j^* & -h_i \end{bmatrix} \begin{bmatrix} y_{ij}(n) \\ y_{ij}^*(n+1) \end{bmatrix} \\ &= \begin{bmatrix} h_i^* & h_j \\ h_j^* & -h_i \end{bmatrix} \begin{bmatrix} v_{ij}(n) \\ v_{ij}^*(n+1) \end{bmatrix} \\ &+ \sqrt{\frac{\alpha}{2}} \left( \begin{bmatrix} |h_i|^2 & h_j h_i^* \\ h_i h_j^* & |h_j|^2 \end{bmatrix} \begin{bmatrix} u_i \\ u_j \end{bmatrix} \\ &+ \begin{bmatrix} |h_j|^2 & -h_j h_i^* \\ -h_i h_j^* & |h_i|^2 \end{bmatrix} \begin{bmatrix} \hat{u}_i \\ \hat{u}_j \end{bmatrix} \right). \end{split}$$

Note that if each node allocates equal power for communicating with its cooperative partner and with the FC, i.e.,  $\alpha = 1/2$ , and there is no error during internode communication, i.e.,  $\hat{u}_i = u_i$ ,  $\hat{u}_j = u_j$ , the given equations reduce to the classical Alamouti's scheme [20]. The new noise terms  $\delta_{ij}^1 =$  $h_i^* v_{ij}(n) + h_j v_{ij}^*(n+1)$  and  $\delta_{ij}^2 = h_j^* v_{ij}(n) - h_i v_{ij}^*(n+1)$ are i.i.d. zero-mean complex Gaussian random variables (RVs) with the variance  $\sigma^2 = (|h_i|^2 + |h_j|^2)\sigma_v^2$ . Next, using the signals  $z_i, z_j$  and the CSI  $h_i, h_j$  for all pairs, the FC forms the LRT  $\Lambda = (f(z_i, z_j \text{ for all pairs} |\mathcal{H}_1)/f(z_i, z_j \text{ for all pairs} |\mathcal{H}_0))$ , where  $f(z_i, z_j \text{ for all pairs} | \mathcal{H}_{\ell})$  indicates the joint pdf of  $z_i, z_j$ corresponding to all pairs given the hypothesis  $\mathcal{H}_{\ell}$ , to make the final decision. In particular, the FC decides on  $\mathcal{H}_1$  when  $\Lambda > \pi_0/\pi_1$  and decides on  $\mathcal{H}_0$  otherwise. We note that  $\mathcal{H}_{\ell} \rightarrow u_i, u_j, \hat{u}_i, \hat{u}_j \rightarrow z_i, z_j$  and  $\mathcal{H}_{\ell} \rightarrow u_i, u_j \rightarrow \hat{u}_i, \hat{u}_j$  form Markov chains. Moreover,  $(z_i, z_j)$  values are independent across the pairs given  $u_i, u_j, \hat{u}_i, \hat{u}_j$  for all pairs. Furthermore, given  $u_i, u_j$ for for all pairs,  $(\hat{u}_i, \hat{u}_j)$  values are independent across the pairs. Therefore, we write

$$f(z_i, z_j \text{ for all pairs} | \mathcal{H}_\ell) = \sum_{u_i} \sum_{u_j} \sum_{\hat{u}_i} \sum_{\hat{u}_j}$$

 $f(z_i, z_j \text{ for all pairs} | u_i, u_j, \hat{u}_i, \hat{u}_j \text{ for all pairs})$ 

$$\times P(u_i, u_j, \hat{u}_i, \hat{u}_j \text{ for all pairs} | \mathcal{H}_{\ell})$$

$$= \sum_{u_i} \sum_{\hat{u}_i} \sum_{u_j} \sum_{\hat{u}_j} \left( \prod_{\text{for all pairs}} f(z_i, z_j | u_i, u_j, \hat{u}_i, \hat{u}_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j)) \right)$$

$$\times P(\hat{u}_i | u_i \text{ for } (\mathcal{S}_i, \mathcal{S}_j)) P(\hat{u}_j | u_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j))$$

$$\times P(u_i, u_j \text{ for all pairs} | \mathcal{H}_{\ell})$$

$$(6)$$

where the sums are taken over all values that  $u_i, u_j, \hat{u}_i, \hat{u}_j$  can assume. Focusing on  $f(z_i, z_j | u_i, u_j, \hat{u}_i, \hat{u}_j$  for  $(S_i, S_j)$ ) in (6), we realize that  $z_i, z_j$  values are conditionally independent complex Gaussian RVs with the variance  $\sigma^2 = (|h_i|^2 + |h_j|^2)\sigma_v^2$ , and the mean values  $\mu_i, \mu_j$  are given in the following equation:

$$\mu_{i} = \sqrt{\frac{\alpha}{2}} \left( |h_{i}|^{2} u_{i} + h_{j} h_{i}^{*} u_{j} + |h_{j}|^{2} \hat{u}_{i} - h_{j} h_{i}^{*} \hat{u}_{j} \right)$$
  
$$\mu_{j} = \sqrt{\frac{\alpha}{2}} \left( h_{i} h_{j}^{*} u_{i} + |h_{j}|^{2} u_{j} - h_{i} h_{j}^{*} \hat{u}_{i} + |h_{i}|^{2} \hat{u}_{j} \right).$$
(7)

Focusing on the term  $P(\hat{u}_i|u_i \text{ for } (S_i, S_j))$  in (6) and considering (4), one can easily verify the following, assuming that the FC only knows the statistics of internode channels  $g_{ij}, g_{ji}$  [23]. Thus

$$P(\hat{u}_i \neq u_i | u_i) = 1 - P(\hat{u}_i = u_i | u_i) = \frac{1}{2} \left( 1 - \sqrt{\frac{\bar{\gamma}_{hs_{ij}}}{1 + \bar{\gamma}_{hs_{ij}}}} \right)$$
  
where  $\bar{\gamma}_{hs_{ij}} = \frac{(1 - \alpha)\sigma_{hs_{ij}}^2}{\sigma_n^2}$  (8)

denotes the average received SNR corresponding to  $(S_i, S_j)$ internode communication. The term  $P(u_i, u_j \text{ for all pairs} | \mathcal{H}_{\ell})$ in (6) can be found in terms of the probability of  $x_i, x_j$ being in certain intervals for all pairs. For instance,  $P(u_i = 1, u_j = -1 \text{ for all pairs} | \mathcal{H}_{\ell}) = P(x_i > \tau_i, x_j < \tau_j \text{ for all pairs} | \mathcal{H}_{\ell})$ . These probabilities can be characterized for the sensing channel model in Section II-A, where  $x_i, x_j$  values are jointly correlated Gaussian RVs with known statistics.<sup>2</sup> When Gaussian sensing noise terms are uncorrelated, we obtain  $P(u_i = 1, u_j = -1 \text{ for all pairs} | \mathcal{H}_0) = \prod_{\{i:u_i=1\}} P_{f_i} \prod_{\{j:u_j=-1\}} (1 - P_{f_j})$  or  $P(u_i = 1, u_j = -1 \text{ for all pairs} | \mathcal{H}_1) = \prod_{\{i:u_i=1\}} P_{d_i} \prod_{\{j:u_j=-1\}} (1 - P_{d_j})$ . When the parameters of sensing channels are unavailable at the FC, the FC cannot apply the LRT. Alternatively, the FC demodulates the channel inputs for all pairs using the signals  $z_i, z_j$  for all pairs and applies the majority rule to the demodulated symbols to reach the final decision.

### IV. COOPERATIVE FUSION ARCHITECTURE WITH SIGNAL FUSION AT SENSORS

# A. Internode Communication Channel Model and Local Decision Rules at Sensors

Suppose nodes  $S_i$  and  $S_j$  are within a pair. Each sensor makes an initial decision based on its measurement. The nodes exchange their decisions over orthogonal channels subject to noise and fading. Let  $u_i$  denote the decision made at  $S_i$  based on  $x_i$ . When  $S_i$  transmits  $\sqrt{1 - \alpha}u_i$ , the received signal at  $S_j$  is  $r_{ij}$ , as shown in (3). Upon receiving  $r_{ij}$ ,  $S_j$  (rather than demodulating the channel input) updates its initial decision by fusing  $r_{ij}$  and its measurement  $x_j$ . In particular,  $S_j$  forms a local LRT  $\tilde{\lambda}_j = f(r_{ij}, x_j | \mathcal{H}_1) / f(r_{ij}, x_j | \mathcal{H}_0)$ , where  $f(r_{ij}, x_j | \mathcal{H}_\ell)$ indicates the joint pdf of  $r_{ij}, x_j$  given the hypothesis  $\mathcal{H}_\ell$ , to make a new decision  $\tilde{u}_j$ . Node  $S_j$  lets  $\tilde{u}_j = 1$  when  $\tilde{\lambda}_j > \pi_0/\pi_1$ and lets  $\tilde{u}_j = -1$  otherwise. Since  $\mathcal{H}_\ell \to u_i \to r_{ij}$  and  $x_j \to$  $\mathcal{H}_\ell, u_i \to r_{ij}$  form Markov chains, we find

$$\tilde{\lambda}_j = \frac{\sum_{u_i} f(r_{ij}|u_i) P(u_i|x_j, \mathcal{H}_1) f(x_j|\mathcal{H}_1)}{\sum_{u_i} f(r_{ij}|u_i) P(u_i|x_j, \mathcal{H}_0) f(x_j|\mathcal{H}_0)}.$$
(9)

Considering  $f(r_{ij}|u_i)$  in (9), we note that given  $g_{ij}$ ,  $u_i$ , we have  $r_{ij} \sim \mathcal{CN}(\sqrt{1-\alpha}u_ig_{ij},\sigma_\eta^2)$ . To find  $P(u_i|x_j,\mathcal{H}_\ell)$  in (9), we note that for the sensing model in Section II-A,  $x_i, x_j$  values are jointly Gaussian RVs with mean  $\ell$ , variances  $\sigma_{w_i}^2, \sigma_{w_j}^2$ , and correlation coefficient  $\rho_{i,j}$  under the hypothesis  $\mathcal{H}_\ell$ . Using the joint pdf of  $x_i, x_j$  and the Bayes' rule  $P(u_i = 1|x_j, \mathcal{H}_\ell) = (P(u_i = 1, x_j | \mathcal{H}_\ell) / f(x_j | \mathcal{H}_\ell))$ , one can show

$$P(u_i = 1 | x_j, \mathcal{H}_\ell) = 1 - P(u_i = -1 | x_j, \mathcal{H}_\ell)$$
$$= Q\left(\frac{\tau_i - \rho_{i,j} x_j \frac{\sigma_{w_i}}{\sigma_{w_j}} - \ell\left(1 - \rho_{i,j} \frac{\sigma_{w_i}}{\sigma_{w_j}}\right)}{\sqrt{\left(1 - \rho_{i,j}^2\right)} \sigma_{w_i}}\right). \quad (10)$$

Finally, we find  $f(x_j|\mathcal{H}_\ell)$  in (9) by noting that given  $\mathcal{H}_\ell$ , we have  $x_j \sim C\mathcal{N}(\ell, \sigma_{w_j}^2)$ . When Gaussian sensing noise terms are uncorrelated,  $u_i$  and  $x_j$  are independent<sup>3</sup> for a given hypothesis  $\mathcal{H}_\ell$ , leading to  $P(u_i|x_j, \mathcal{H}_\ell) = P(u_i|\mathcal{H}_\ell)$ . The pair  $(\mathcal{S}_i, \mathcal{S}_j)$  sends  $\sqrt{\alpha}\tilde{u}_i, \sqrt{\alpha}\tilde{u}_j$  to the FC over two orthogonal channels subject to noise and fading. Considering the definitions of channel variances, we note that, effectively,  $\mathcal{S}_i$  spends  $(1 - \alpha)\mathcal{P}$  and  $\alpha\mathcal{P}$ , respectively, for internode and sensor-FC communication.

<sup>&</sup>lt;sup>2</sup>For non-Gaussian  $w_k$ 's, the probability  $P(u_i, u_j \text{ for all pairs} | \mathcal{H}_\ell)$  should be calculated in terms of the joint pdf  $f(w_1, \ldots, w_K)$ .

<sup>&</sup>lt;sup>3</sup>For non-Gaussian sensing noise terms, when forming  $\tilde{\lambda}_j$ ,  $P(u_i|x_j, \mathcal{H}_\ell)$ and  $f(x_j|\mathcal{H}_\ell)$  should be calculated, respectively, based on the joint pdf  $f(w_i, w_j)$  and the pdf  $f(w_j)$ . Moreover,  $P(u_i|x_j, \mathcal{H}_\ell) = P(u_i|\mathcal{H}_\ell)$  when  $w_i, w_j$  values are mutually independent.

# *B.* Node–FC Communication Channel Model and Fusion Rule at the FC

Let  $y_i$  and  $y_j$  denote the received signals at the FC corresponding to the pair  $(S_i, S_j)$ . We have

$$\begin{split} y_i &= \sqrt{\alpha} \tilde{u}_i h_i + v_i, \ y_j &= \sqrt{\alpha} \tilde{u}_j h_j + v_j \\ h_{i(j)} &\sim \mathcal{CN}\left(0, \sigma_{h_{i(j)}}^2\right), v_i, \ v_j &\sim \mathcal{CN}\left(0, \sigma_v^2\right). \end{split}$$

The terms  $v_i$  and  $v_j$  are the receiver noise at the FC. We assume that noise and fading coefficients are independent, and noise terms are i.i.d. across the pairs. Next, using the signals  $y_i, y_j$  and the CSI  $h_i, h_j$  for all pairs, the FC forms the LRT  $\Lambda = (f(y_i, y_j \text{ for all pairs} | \mathcal{H}_1) / f(y_i, y_j \text{ for all pairs} | \mathcal{H}_0))$ , where  $f(y_i, y_j \text{ for all pairs} | \mathcal{H}_\ell)$  indicates the joint pdf of  $y_i, y_j$ corresponding to all pairs given the hypothesis  $\mathcal{H}_\ell$ , to make the final decision. We note that  $\mathcal{H}_\ell \to \tilde{u}_i, \tilde{u}_j \to y_i, y_j$  form a Markov chain. Furthermore,  $(y_i, y_j)$  values are independent across the pairs given  $\tilde{u}_i, \tilde{u}_j$  for all pairs. Hence, we can write

$$f(y_i, y_j \text{ for all pairs} | \mathcal{H}_{\ell}) = \sum_{\tilde{u}_i} \sum_{\tilde{u}_j} \left( P(\tilde{u}_i, \tilde{u}_j \text{ for all pairs} | \mathcal{H}_{\ell}) \times f(y_i, y_j \text{ for all pairs} | \tilde{u}_i, \tilde{u}_j \text{ for all pairs} ) \right)$$
$$= \sum_{\tilde{u}_i} \sum_{\tilde{u}_j} \left( \prod_{\text{for all pairs}} f(y_i, y_j | \tilde{u}_i, \tilde{u}_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j)) \right) \times P(\tilde{u}_i, \tilde{u}_j \text{ for all pairs} | \mathcal{H}_{\ell})$$
(11)

where the sums are taken over all values that  $\tilde{u}_i, \tilde{u}_j$  can take. To find  $f(y_i, y_j | \tilde{u}_i, \tilde{u}_j$  for  $(S_i, S_j))$  in (11), we realize that in the pair  $(S_i, S_j)$  given  $\tilde{u}_i, \tilde{u}_j$ , the variables  $y_i$ and  $y_j$  are independent complex Gaussian RVs with mean values  $\mu_i = \sqrt{\alpha} \tilde{u}_i h_i$  and  $\mu_j = \sqrt{\alpha} \tilde{u}_j h_j$  and variance  $\sigma_v^2$ . To obtain  $P(\tilde{u}_i, \tilde{u}_j$  for all pairs $|\mathcal{H}_\ell)$  in (11), we assume that the FC only knows the statistics of internode channels  $g_{ij}, g_{ji}$ . Note that  $\mathcal{H}_\ell \to x_i, x_j, r_{ij}, r_{ji} \to \tilde{u}_i, \tilde{u}_j$  forms a Markov chain. Therefore

$$P(\tilde{u}_i, \tilde{u}_j \text{ for all pairs} | \mathcal{H}_{\ell}) = \int_{g_{ij}} \int_{g_{ji}} \int_{x_i} \int_{x_j} \int_{r_{ij}} \int_{r_{ji}} \int_{r_{ji}} \int_{r_{ji}} P(\tilde{u}_i, \tilde{u}_j \text{ for all pairs} | g_{ij}, g_{ji}, x_i, x_j, r_{ij}, r_{ji} \text{ for all pairs}) \times f(x_i, x_j, r_{ij}, r_{ji} \text{ for all pairs} | \mathcal{H}_{\ell}) \times f(g_{ij}, g_{ji}) dg_{ij} dg_{ji} dx_i dx_j dr_{ij} dr_{ji}.$$
(12)

Equation (12) can be further simplified by noting that given  $x_i, x_j, r_{ij}, r_{ji}$  for all pairs, the variables  $(\tilde{u}_i, \tilde{u}_j)$  are inde-

pendent across the pairs. Furthermore, within a pair, given  $x_i, x_j, r_{ij}, r_{ji}$ , variables  $\tilde{u}_i$  and  $\tilde{u}_j$  are independent. In addition, we note that  $\mathcal{H}_{\ell} \to x_i, x_j \to r_{ij}, r_{ji}$  forms a Markov chain, and given  $x_i, x_j$  for all pairs,  $(r_{ij}, r_{ji})$  values are independent across the pairs. Moreover, within a pair, given  $x_i, x_j$ , variables  $r_{ij}, r_{ji}$  are independent. Combining all, we can rewrite (12) as (13), shown at the bottom of the page.

Focusing on the term  $P(\tilde{u}_i|x_i, r_{ji} \text{ for } (\mathcal{S}_i, \mathcal{S}_j))$  in (13), we note that  $P(\tilde{u}_i = 1|x_i, r_{ji}) = \mathbf{1}_{\{\tilde{\lambda}_i > \pi_0/\pi_1\}}$  and  $P(\tilde{u}_i = -1|x_i, r_{ji}) = \mathbf{1}_{\{\tilde{\lambda}_i < \pi_0/\pi_1\}}$ , where  $\tilde{\lambda}_i$  depends on internode channels  $g_{ji}$ , threshold  $\tau_j$ , sensing noise variances  $\sigma_{w_i}^2, \sigma_{w_j}^2$ , and correlation coefficient  $\rho_{i,j}$ . Similarly, we can find  $P(\tilde{u}_j|x_j, r_{ij} \text{ for } (\mathcal{S}_i, \mathcal{S}_j))$  in (13). Considering the term  $f(r_{ji}|x_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j))$  in (13), we find

$$f(r_{ji}|x_j \text{ for } (S_i, S_j))$$
  
=  $f(r_{ji}|u_j = 1)P(u_j = 1|x_j) + f(r_{ji}|u_j = -1)P(u_j = -1|x_j)$   
=  $f(r_{ji}|u_j = 1)\mathbf{1}_{\{x_j > \tau_j\}} + f(r_{ji}|u_j = -1)\mathbf{1}_{\{x_j < \tau_j\}}$  (14)

where  $f(r_{ji}|u_j)$  in (14) can be found noting that given  $g_{ji}, u_j$ , we have  $r_{ji} \sim C\mathcal{N}(\sqrt{1-\alpha}u_jg_{ji}, \sigma_\eta^2)$ . Considering the term  $f(x_i, x_j)$  for all pairs $|\mathcal{H}_\ell|$  in (13), we note that when Gaussian sensing noise terms are uncorrelated<sup>4</sup> we obtain  $f(x_i, x_j)$  for all pairs $|\mathcal{H}_\ell| = \prod_{\text{for all pairs}} f(x_i \text{ for } (\mathcal{S}_i, \mathcal{S}_j)|\mathcal{H}_\ell)$  $f(x_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j)|\mathcal{H}_\ell)$ . Combining all these, one can verify that the probability in (12) depends on sensing channels through  $\tau_i, \tau_j, \sigma_{w_i}^2, \sigma_{w_j}^2, \rho_{i,j}$  and the average received SNR  $\bar{\gamma}_{hs_{ij}}$  corresponding to  $(\mathcal{S}_i, \mathcal{S}_j)$  internode communication. When the parameters of sensing channels are unavailable at the FC, the FC demodulates the channel inputs for all pairs using the signals  $y_i, y_j$  for all pairs and applies the majority rule to the demodulated symbols to reach the final decision.

### V. PARALLEL FUSION ARCHITECTURE WITH LOCAL THRESHOLD CHANGING AT SENSORS

#### A. Local Decision Rules at Sensors

Suppose nodes  $S_i$  and  $S_j$  are within a pair. Each sensor makes an initial decision based on its measurement. Let  $u_i$ denote the decision made at  $S_i$  based on  $x_i$ . In the absence of internode communication,  $S_i$  assumes that the decision  $u_j$ (made at  $S_j$  based on  $x_j$ ) is different from  $u_i$ , i.e.,  $S_i$  assumes  $u_j = -u_i$ . Next,  $S_i$  forms another decision  $\bar{u}_i$  by fusing the

$$\int_{g_{ij}g_{ji}} \int_{x_i} \int_{x_j} \int_{r_{ij}r_{ji}} \left( \prod_{\text{for all pairs}} P\left(\tilde{u}_i | x_i, r_{ji} \text{ for } (\mathcal{S}_i, \mathcal{S}_j)\right) P\left(\tilde{u}_j | x_j, r_{ij} \text{ for } (\mathcal{S}_i, \mathcal{S}_j)\right) \times f\left(r_{ji} | x_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j)\right) f\left(r_{ij} | x_i \text{ for } (\mathcal{S}_i, \mathcal{S}_j)\right) \right)$$

 $\times f(x_i, x_j \text{ for all pairs } |\mathcal{H}_\ell) f(g_{ij}, g_{ji}) dg_{ij} dg_{ji} dx_i dx_j dr_{ij} dr_{ji}$ (13)

<sup>&</sup>lt;sup>4</sup>For non-Gaussian sensing noise,  $\overline{\lambda}_i$  would change, as explained in the previous footnote. Moreover,  $P(u_j|x_j)$  in (14) and  $f(x_i, x_j)$  for all pairs  $|\mathcal{H}_\ell|$  in (13), respectively, should be calculated based on the pdf  $f(w_j)$  and the joint pdf  $f(w_1, \ldots, w_K)$ .

assumed decision  $u_j$  and its measurement  $x_i$ . In particular,  $S_i$  forms a local LRT  $\bar{\lambda}_i = f(x_i, u_j = -u_i | \mathcal{H}_1) / f(x_i, u_j = -u_i | \mathcal{H}_0)$ , where  $f(x_i, u_j = -u_i | \mathcal{H}_\ell)$  indicates the joint pdf of  $x_i$  and the assumed decision  $u_j$  at  $S_i$  given the hypothesis  $\mathcal{H}_\ell$ , to make  $\bar{u}_i$ . Node  $S_i$  lets  $\bar{u}_i = 1$  when  $\bar{\lambda}_i > \pi_0 / \pi_1$  and lets  $\bar{u}_i = -1$  otherwise. We find<sup>5</sup>

$$\bar{\lambda}_i = \frac{P(u_j = -u_i | x_i, \mathcal{H}_1) f(x_i | \mathcal{H}_1)}{P(u_j = -u_i | x_i, \mathcal{H}_0) f(x_i | \mathcal{H}_0)}$$
(15)

in which  $P(u_j = -u_i | x_i, \mathcal{H}_\ell)$  is given in (10), and  $f(x_j | \mathcal{H}_\ell)$  is found by noting<sup>6</sup> that given  $\mathcal{H}_\ell$ , we have  $x_j \sim C\mathcal{N}(\ell, \sigma_{w_j}^2)$ . In fact, one can verify that  $S_i$  finds  $u_i, \bar{u}_i$ , as follows:

$$\begin{cases}
 u_i = 1, \bar{u}_i = 1, & \text{if } x_i > \tau'_{i_1} \\
 u_i = -1, \bar{u}_i = -1, & \text{if } x_i < \tau'_{i_2} \\
 u_i = -1, \bar{u}_i = 1, & \text{if } \tau'_{i_2} < x_i < \tau_i \\
 u_i = 1, \bar{u}_i = -1, & \text{if } \tau_i < x_i < \tau'_{i_1}
\end{cases}$$
(16)

where the thresholds  $\tau'_{i_1}, \tau'_{i_2}$ , given in the footnote, depend on  $\sigma^2_{w_i}, \rho_{i,j}$  and satisfy  $\tau'_{i_2} < \tau_i < \tau'_{i_1}$ . When Gaussian sensing noise terms are uncorrelated, the assumed decision  $u_j$ and  $x_i$  are independent for a given hypothesis  $\mathcal{H}_{\ell}$ , leading to  $P(u_j = -u_i | x_i, \mathcal{H}_{\ell}) = P(u_j = -u_i | \mathcal{H}_{\ell})$  in (15). Consequently, the local LRT  $\lambda_i$  in (15) can be further simplified, and  $\tau'_{i_1}, \tau'_{i_2}$  in (16), respectively, reduce to  $\tau_{i_1} = 0.5 + \sigma^2_{w_i} \ln((1 - P_{d_j})\pi_1)$  and  $\tau_{i_2} = 0.5 + \sigma^2_{w_i} \ln(P_{f_j}\pi_0/P_{d_j}\pi_1)$ , where  $\tau_{i_2} < \tau_i < \tau_{i_1}$ . Since, in addition to threshold  $\tau_i$  employed in scheme 1),  $x_i$  is also compared against two additional thresholds  $\tau'_{i_1}, \tau'_{i_2}$ , we refer to scheme 3) as "local threshold

<sup>5</sup>Consider a hypothetical case where  $S_i$  assumes  $u_j = -1$  and makes a decision  $u_i^0$  by optimally fusing the assumed  $u_j = -1$  and  $x_i$ . In particular,  $S_i$  lets  $u_i^0 = 1$  when  $\lambda_i^0 = (f(x_i, u_j = -1|\mathcal{H}_1)/f(x_i, u_j = -1|\mathcal{H}_0)) > (\pi_0/\pi_1)$  and lets  $u_i^0 = -1$  otherwise. Consider another hypothetical case where  $S_i$  assumes  $u_j = 1$  and makes a decision  $u_i^1$  by optimally fusing the assumed  $u_j = 1$  and  $x_i$ . In particular,  $S_i$  lets  $u_i^1 = 1$  when  $\lambda_i^1 = (f(x_i, u_j = 1|\mathcal{H}_0)) > (\pi_0/\pi_1)$  and lets  $u_i^1 = -1$  otherwise. One can verify  $u_i^0 = 1$  when  $x_i > \tau'_{i_1}$  and  $u_i^0 = -1$  otherwise, also  $u_i^1 = 1$  when  $x_i > \tau'_{i_2}$  and  $u_i^1 = -1$  otherwise, where  $\tau'_{i_1} = 0.5 + \sigma_{w_i}^2 \ln(P(u_j = -1|x_i, \mathcal{H}_0)\pi_0/P(u_j = -1|x_i, \mathcal{H}_1)\pi_1)$  and  $\tau'_{i_2} = 0.5 + \sigma_{w_i}^2 \ln(P(u_j = 1|x_i, \mathcal{H}_0)\pi_0/P(u_j = 1|x_i, \mathcal{H}_1)\pi_1)$  and  $\tau'_{i_2} < \tau_i < \tau'_{i_1}$ . For these hypothetical cases, now suppose  $x_i < \tau_i$ , and thus,  $u_i = -1$ . Since  $\tau_i < \tau'_{i_1}$ , we have  $x_i < \tau'_{i_1}$ , and thus,  $u_i^0 = -1$ , i.e.,  $u_i^0 = u_i$ , wherea  $u_i^1 \ ant \pm 1$ . Therefore, the useful information is embedded in  $u_i, u_i^1$ , as  $u_i^0$  conveys no extra information. Similarly, one can argue that when  $x_i > \tau_i$ , the useful information is one solve yextra information. In conclusion, node  $S_i$  should assume  $u_j = -u_i$  to be able to extract more information from  $x_i$ .

<sup>6</sup>For non-Gaussian sensing noise, similar to  $\overline{\lambda}_i$  in Section IV-A,  $\overline{\lambda}_i$  would change. In particular,  $P(u_j|x_i, \mathcal{H}_\ell)$  and  $f(x_i|\mathcal{H}_\ell)$  should be calculated, respectively, based on the joint pdf  $f(w_i, w_j)$  and the pdf  $f(w_i)$ .

changing." The pair  $(S_i, S_j)$  sends  $u_i, \bar{u}_i, u_j, \bar{u}_j$  to the FC in two consecutive time slots, exploiting Alamouti's STC scheme. In particular, in the *n*th slot,  $S_i$  and  $S_j$  can simultaneously send  $(u_i/\sqrt{2})$  and  $(\bar{u}_j/\sqrt{2})$ , respectively. In the (n + 1)th time slot,  $S_i$  and  $S_j$  can simultaneously send  $-(\bar{u}_i/\sqrt{2})$  and  $(u_j/\sqrt{2})$ , respectively. Considering the definitions of channel variances, we note that, effectively,  $S_i$  spends  $\mathcal{P} = (1/2)\mathcal{P} + (1/2)\mathcal{P}$  for sensor–FC communication.

# *B.* Node–FC Communication Channel Model and Fusion Rule at FC

Let  $y_{ij}(n)$  and  $y_{ij}(n + 1)$  denote the received signals at the FC corresponding to the pair  $(S_i, S_j)$  during two consecutive time slots. The signal model in (5) still holds true, after substituting  $\sqrt{\alpha/2}$  with  $1/\sqrt{2}$ ,  $u_j$  with  $\bar{u}_j$ ,  $\hat{u}_j$  with  $\bar{u}_i$ , and  $\hat{u}_i$  with  $u_j$ . We can take similar steps as in Section III-B to find the signals  $z_i, z_j$  using  $y_{ij}(n), y_{ij}^*(n + 1)$ . Next, using the signals  $z_i, z_j$  and the CSI  $h_i, h_j$  for all pairs, the FC forms the LRT  $\Lambda = (f(z_i, z_j \text{ for all pairs} |\mathcal{H}_1) / f(z_i, z_j \text{ for all pairs} |\mathcal{H}_0))$ . We note that  $\mathcal{H}_\ell \to u_i, u_j, \bar{u}_i, \bar{u}_j \to z_i, z_j$  forms a Markov chain. Moreover,  $(z_i, z_j)$  values are independent across the pairs given  $u_i, u_j, \bar{u}_i, \bar{u}_j$  for all pairs. Therefore, we reach (17), shown at the bottom of the page.

Focusing on the term  $f(z_i, z_j | u_i, u_j, \bar{u}_i, \bar{u}_j$  for  $(\mathcal{S}_i, \mathcal{S}_j)$  in (17), we realize that  $z_i, z_j$  values are conditionally independent complex Gaussian RVs with variance  $\sigma^2 = (|h_i|^2 + |h_j|^2)\sigma_v^2$ and mean values  $\mu_i, \mu_j$  given in (7), after substituting  $\sqrt{(\alpha/2)}$ with  $(1/\sqrt{2})$ ,  $u_j$  with  $\bar{u}_j$ ,  $\hat{u}_j$  with  $\bar{u}_i$ , and  $\hat{u}_i$  with  $u_j$ . To find  $P(u_i, u_j, \bar{u}_i, \bar{u}_j)$  for all pairs  $|\mathcal{H}_\ell)$ , we note that this term can be expressed in terms of the probability of  $x_i, x_j$  being in certain intervals for all pairs. For instance,  $P(u_i =$  $\bar{u}_i = 1, u_j = \bar{u}_j = -1$  for all pairs $|\mathcal{H}_\ell) = P(x_i > \tau_{ij_1}, x_j < 1)$  $\tau_{ij_2}$  for all pairs  $|\mathcal{H}_{\ell}|$ . Now, these probabilities can be easily characterized for the sensing channel model in Section II-A, where  $x_i, x_j$  values are jointly correlated Gaussian RVs with known statistics.7 When Gaussian sensing noise terms are uncorrelated, the decisions  $(u_i, \bar{u}_i)$  and  $(u_i, \bar{u}_i)$  given hypothesis  $\mathcal{H}_{\ell}$  are independent across the pairs. Therefore,  $P(u_i, u_j, u_j)$  $\bar{u}_i, \bar{u}_j$  for all pairs  $|\mathcal{H}_\ell| = \prod_{\text{for all pairs}} P(u_i, \bar{u}_i \text{ for } (\mathcal{S}_i, \mathcal{S}_j) | \mathcal{H}_\ell) \times$  $P(u_j, \bar{u}_j \text{ for } (\mathcal{S}_i, \mathcal{S}_j) | \mathcal{H}_\ell)$ . When the parameters of sensing channels are unavailable at the FC, the FC demodulates the channel inputs for all pairs using the signals  $z_i, z_j$  for all pairs and applies the majority rule to the demodulated symbols to reach the final decision.

<sup>7</sup>For non-Gaussian  $w_k$ 's,  $P(u_i, u_j, \bar{u}_i, \bar{u}_j \text{ for all pairs} | \mathcal{H}_\ell)$  should be calculated in terms of the joint pdf  $f(w_1, \ldots, w_K)$ .

$$f(z_i, z_j \text{ for all pairs} | \mathcal{H}_{\ell}) = \sum_{u_i \text{ for all pairs } u_j \text{ for all pairs } \bar{u}_i \text{ for all pairs } \bar{u}_i \text{ for all pairs}} \sum_{\bar{u}_i \text{ for all pairs}} \sum_{\bar{u}_j \text{ fo$$

#### VI. PERFORMANCE ANALYSIS

Section VI-A provides an upper bound on the average error probability for scheme 1) in Section III. Leveraging on this, Section VI-B and C, respectively, provide upper bounds on the average error probability for the classical parallel fusion architecture in Section II-B and scheme 2) in Section IV. For mathematical tractability, we assume that the Gaussian sensing noise terms  $w_k$  are identically distributed and uncorrelated,<sup>8</sup> i.e., we have  $\sigma_{w_k}^2 = \sigma_w^2$ , and  $\tau_k = \tau$ , and thus,  $P_{d_k} = P_d$ , and  $P_{f_k} = P_f$ . Moreover, we assume that sensors are positioned equally distant from the FC, and thus,  $\bar{\gamma}_h^2 = (\sigma_h^2/\sigma_v^2)$ . Furthermore, distances between the cooperative partners are assumed equal across the pairs, and therefore,  $\bar{\gamma}_{hs}^2 = ((1 - \alpha)\sigma_{hs}^2/\sigma_\eta^2)$ . In Section VII, we validate the analytical results of this section via Monte Carlo simulations.

### A. Cooperative Fusion Architecture With STC at Sensors

For performance analysis, suppose each pair of cooperative partners  $(S_i, S_j)$  is associated with a unique group index s, where  $s = 1, \ldots, S$  and S = K/2. We denote the two nodes within group s as  $(S_{2s-1}, S_{2s})$ , i.e., we map the indexes i, j in Section III into 2s - 1, 2s, respectively. For the LRT fusion rule in Section III-B, the conditional error probability is  $P_{e|h} = P_{e_1|h} + P_{e_2|h}$ , where  $P_{e_1|h} = P(\Lambda > (\pi_0/\pi_1)|\mathcal{H}_0)\pi_0$  and  $P_{e_2|h} = P(\Lambda < (\pi_0/\pi_1)|\mathcal{H}_1)\pi_1$  and  $\Lambda = (f(z_{2s-1}, z_{2s} \text{ for } s = 1, \ldots, S|\mathcal{H}_1)/f(z_{2s-1}, z_{2s} \text{ for } s = 1, \ldots, S|\mathcal{H}_1)/f(z_{2s-1}, z_{2s} \text{ for } s = 1, \ldots, S|\mathcal{H}_0)$ , conditioned on the channel coefficients  $h_{2s-1}, h_{2s}$  for  $s = 1, \ldots, S$  at the FC. The average error probability  $\bar{P}_e = \bar{P}_{e_1} + \bar{P}_{e_2}$  is obtained by taking the averages of  $P_{e_1|h}, P_{e_2|h}$  over the distribution of the channel coefficients. Our goal is to provide upper bounds on  $P_{e_1|h}, P_{e_2|h}$  and their corresponding averages  $\bar{P}_{e_1} = \mathbb{E}\{P_{e_1|h}\}, \bar{P}_{e_2} = \mathbb{E}\{P_{e_2|h}\}$ . We use the following notation in our derivations. To

<sup>8</sup>The derivations in this section also hold true for i.i.d. sensing noise terms.

capture all different values that  $u_{2s-1}, u_{2s}, \hat{u}_{2s-1}, \hat{u}_{2s}$  for  $s=1,\ldots,S$  can take, we consider two K-length sequences  $(a_{n_1}^1,\ldots,a_{n_1}^{2s-1},a_{n_1}^{2s},\ldots,a_{n_1}^{2S})$  and  $(a_{m_1}^1,\ldots,a_{m_1}^{2s-1},a_{m_1}^{2s},\ldots,a_{m_1}^{2S})$ , where  $a_{n_1}^{2s-1},a_{n_1}^{2s}\in\{1,-1\}$  are the values assumed by  $u_{2s-1},u_{2s}$ , and  $a_{m_1}^{2s-1},a_{m_1}^{2s}\in\{1,-1\}$  are the values assumed by  $\hat{u}_{2s-1},\hat{u}_{2s}$ . Moreover, let  $n_1$  and  $m_1$ , respectively, be the decimal numbers corresponding to the two K-length binary sequences, when those  $a_{n_1}^k$  and  $a_{m_1}^k$  assuming a -1 value in the sequences are reassigned a 0 value. Let  $Q_{n_1}$  denote the number of 1's in the sequence  $(a_{n_1}^1,a_{n_1}^2,\ldots,a_{n_1}^{2s-1},a_{n_1}^{2s})$ . Define  $F_{n_1,m_1},F_{n_1},T_{n_1,m_1},d_{n_1,m_1}$  as (18)–(21), shown at the bottom of the page.

Furthermore, let M be an integer that satisfies  $\begin{array}{l} (P_d^M(1-P_d)^{K-M}/P_f^M(1-P_f)^{K-M}) > (\pi_0/\pi_1) & \text{and} \\ (P_d^{M-1}(1-P_d)^{K-M+1}/P_f^{M-1}(1-P_f)^{K-M+1}) < (\pi_0/\pi_1), \end{array}$ i.e., M is the smallest number of nodes that can decide  $\mathcal{H}_1$  while  $\mathcal{H}_1$  is true and the FC decides correctly, had all communication channels been error free. Define the sets  $S_0 = \{d_{n_1,m_1} \text{ where } Q_{n_1} < M\}$  and  $S_1 = \{d_{n_1,m_1} \text{ where }$  $Q_{n_1} \geq M$ . Note that all entries of  $S_0$  and  $S_1$  are, respectively, negative and positive. Let  $|S_0|$  and  $|S_1|$  denote the cardinalities of  $S_0$  and  $S_1$ , respectively. Utilizing the given notations, we can rewrite  $f(z_{2s-1}, z_{2s} \text{ for } s = 1, \dots, S | \mathcal{H}_{\ell})$  in  $\Lambda$  as (22), shown at the bottom of the page. Since sensing noise terms are identically distributed and uncorrelated, we find  $P(F_{n_1}|\mathcal{H}_\ell) = \prod_{s=1}^{S} P(u_{2s-1} = a_{n_1}^{2s-1} | \mathcal{H}_\ell) P(u_{2s} = a_{n_1}^{2s} | \mathcal{H}_\ell)$  in (22), and thus,  $P(F_{n_1}|\mathcal{H}_1) = P_d^{Q_{n_1}} (1 - P_d)^{K - Q_{n_1}}$ , and  $P(F_{n_1}|\mathcal{H}_0) = P_f^{Q_{n_1}}(1-P_f)^{K-Q_{n_1}}$ . The term  $T_{n_1,m_1}$  in (22) is calculated using (8) and depends on the average received SNR  $\bar{\gamma}_{hs}$  corresponding to internode communication. Combining (21) and (22), we find (23) and (24), shown at the bottom of the next page, where  $\mathcal{T}_{e_1|h} = P(\sum_{n_1,m_1} d_{n_1,m_1} >$  $0|F_{n,m})$  in (23) and  $\mathcal{T}_{e_2|h} = 1 - \mathcal{T}_{e_1|h}$  in (24). Note that  $P_{e_1|h}$  in (23) and  $P_{e_2|h}$  in (24) depend on  $h_{2s-1}, h_{2s}$  for  $s = 1, \ldots, S$  only through  $\mathcal{T}_{e_1|h}$  and  $\mathcal{T}_{e_2|h}$ , respectively. These imply that the problem of finding upper bounds on  $P_{e_1|h}$ ,  $P_{e_2|h}$ 

$$F_{n_1,m_1} = \left\{ u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}, \hat{u}_{2s-1} = a_{m_1}^{2s-1}, \hat{u}_{2s} = a_{m_1}^{2s} \text{ for } s = 1, \dots, S \right\}$$
(18)

$$F_{n_1} = \left\{ u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s} \text{ for } s = 1, \dots, S \right\}$$
(19)

$$T_{n_1,m_1} = \prod_{s=1}^{n} P\left(\hat{u}_{2s-1} = a_{m_1}^{2s-1} | u_{2s-1} = a_{n_1}^{2s-1}\right) P\left(\hat{u}_{2s} = a_{m_1}^{2s} | u_{2s} = a_{n_1}^{2s}\right)$$
(20)

$$d_{n_1,m_1} = \left(\pi_1 P_d^{Q_{n_1}} (1 - P_d)^{K - Q_{n_1}} - \pi_0 P_f^{Q_{n_1}} (1 - P_f)^{K - Q_{n_1}}\right) \\ \times T_{n_1,m_1} \prod_{s=1}^S f\left(z_{2s-1}, z_{2s} | u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}, \hat{u}_{2s-1} = a_{m_1}^{2s-1}, \hat{u}_{2s} = a_{m_1}^{2s}\right)$$
(21)

 $f(z_{2s-1}, z_{2s} \text{ for } s = 1, \dots, S | \mathcal{H}_{\ell})$ 

$$=\sum_{n_1,m_1} P(F_{n_1}|\mathcal{H}_\ell) T_{n_1,m_1} \prod_{s=1}^S f\left(z_{2s-1}, z_{2s}|u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}, \hat{u}_{2s-1} = a_{m_1}^{2s-1}, \hat{u}_{2s} = a_{m_1}^{2s}\right) \quad (22)$$

and their corresponding averages  $\bar{P}_{e_1}$ ,  $\bar{P}_{e_2}$  can be reduced to finding upper bounds on  $\mathcal{T}_{e_1|h}$ ,  $\mathcal{T}_{e_2|h}$  and their respective averages  $\bar{\mathcal{T}}_{e_1} = \mathbb{E}\{\mathcal{T}_{e_1|h}\}, \bar{\mathcal{T}}_{e_2} = \mathbb{E}\{\mathcal{T}_{e_2|h}\}$ . In the Appendix, we establish (25)–(29), shown at the bottom of the page.

The upper bound on  $\overline{\mathcal{T}}_{e_2}$  depends on t. Our simulations indicate that the bound is minimized for  $t \approx 0.3$ . Furthermore, the upper bounds depend on the power allocation parameter  $\alpha$ . In Section VII, we investigate the optimal  $\alpha$  that minimizes these bounds. In the ideal case when the internode communication is error free, we find  $a_n^{2s-1} = a_m^{2s-1}$  and  $a_n^{2s} = a_m^{2s}$ , and thus, (28) and (29), respectively, reduce to  $\mathcal{D}_1(n, m, n_1, m_1) = (1 + (\alpha \overline{\gamma}_h \overline{a}_1/8))^{-2}$  and  $\mathcal{D}_2(n, m, n_1, m_1) = (1 + (\alpha (t^2 - t) \overline{\gamma}_h \overline{a}_1/2))^{-2}$ .

communication renders the notations in Section VI-A simple, as the indexes  $m, m_1$  and the decisions  $\hat{u}_{2s-1}, \hat{u}_{2s}$  are dropped,  $z_{2s-1}, z_{2s}$  are substituted with  $y_{2s-1}, y_{2s}$ , and the noise terms  $\delta_s^1, \delta_s^2$  are substituted with  $v_{2s-1}, v_{2s}$ . Consequently, the derivations of the upper bounds become rather easy. In particular, instead of  $d_{n_1,m_1}$  in (21), we define  $d_{n_1}$  in (30), shown at the bottom of the page. Moreover, the relationship between  $P_{e_1|h}, \mathcal{T}_{e_1|h}$  in (23) and  $P_{e_2|h}, \mathcal{T}_{e_2|h}$  in (24) can be revised as follows:

$$P_{e_1|h} = \pi_0 P\left(\sum_{n_1} d_{n_1} > 0|\mathcal{H}_0\right) = \pi_0 \sum_n \mathcal{T}_{e_1|h} P(F_n|\mathcal{H}_0)$$
$$P_{e_2|h} = \pi_1 P\left(\sum_{n_1} d_{n_1} < 0|\mathcal{H}_1\right) = \pi_1 \sum_n \mathcal{T}_{e_2|h} P(F_n|\mathcal{H}_1)$$
(31)

#### B. Classical Parallel Fusion Architecture

Leveraging on the derivations in Section VI-A, we provide upper bounds on  $\overline{\mathcal{T}}_{e_1}, \overline{\mathcal{T}}_{e_2}$ . In fact, the absence of internode

in which  $F_n$  is defined in (19). We redefine  $S_0 = \{d_{n_1}, \text{ where } Q_{n_1} < M\}$  and  $S_1 = \{d_{n_1}, \text{ where } Q_{n_1} \ge M\}$ ,

$$P_{e_1|h} = \pi_0 P\left(\sum_{n_1,m_1} d_{n_1,m_1} > 0|\mathcal{H}_0\right) = \pi_0 \sum_{n,m} \mathcal{T}_{e_1|h} P(F_{n,m}|\mathcal{H}_0) = \pi_0 \sum_{n,m} \mathcal{T}_{e_1|h} P_f^{Q_n} (1 - P_f)^{K - Q_n} T_{n,m}$$
(23)

$$P_{e_2|h} = \pi_1 P\left(\sum_{n_1,m_1} d_{n_1,m_1} < 0|\mathcal{H}_1\right) = \pi_1 \sum_{n,m} \mathcal{T}_{e_2|h} P(F_{n,m}|\mathcal{H}_1) = \pi_1 \sum_{n,m} \mathcal{T}_{e_2|h} P_d^{Q_n} (1 - P_d)^{K - Q_n} T_{n,m}$$
(24)

$$\bar{\mathcal{T}}_{e_1} < \frac{\mathbf{1}_{\{Q_n < M\}}}{2\sqrt{|S_1|}} \sum_{d_{n_1,m_1} \in S_1} \left[ \sqrt{G(n,m,n_1,m_1)} \prod_{s=1}^S \mathcal{D}_1(n,m,n_1,m_1) \right] + \mathbf{1}_{\{Q_n \ge M\}}$$
(25)

$$\bar{\mathcal{T}}_{e_2} < \frac{\mathbf{1}_{\{Q_n > M\}}}{|S_0|} \sum_{d_{n_1,m_1} \in S_0} \left[ \min_t \left( |S_0| G(n,m,n_1,m_1) \right)^t \prod_{s=1}^S \mathcal{D}_2(n,m,n_1,m_1) \right] + \mathbf{1}_{\{Q_n \le M\}}$$
(26)

$$G(n,m,n_1,m_1) = \frac{T_{n_1,m_1}}{T_{n,m}} \times \frac{\left(\pi_1 P_d^{Q_{n_1}} (1-P_d)^{K-Q_{n_1}} - \pi_0 P_f^{Q_{n_1}} (1-P_f)^{K-Q_{n_1}}\right)}{\left(\pi_0 P_f^{Q_n} (1-P_f)^{K-Q_n} - \pi_1 P_d^{Q_n} (1-P_d)^{K-Q_n}\right)}$$
(27)

$$\mathcal{D}_{1}(n,m,n_{1},m_{1}) = \left( \left( 1 + \frac{\alpha \bar{\gamma}_{h} \bar{a}_{1}}{8} \right) \left( 1 + \frac{\alpha \bar{\gamma}_{h} \bar{a}_{2}}{8} \right) - \frac{\alpha^{2} \bar{\gamma}_{h}^{2} \bar{a}_{3}}{64} \right)^{-1}$$

$$\mathcal{D}_{2}(n,m,n_{1},m_{1}) = \left( \left( 1 + \frac{\alpha (t^{2} - t) \bar{\gamma}_{h} \bar{a}_{1}}{2} \right) \left( 1 + \frac{\alpha (t^{2} - t) \bar{\gamma}_{h} \bar{a}_{2}}{2} \right) - \frac{\alpha^{2} (t^{2} - t)^{2} \bar{\gamma}_{h}^{2} \bar{a}_{3}}{16} \right)^{-1}$$

$$\bar{a}_{1} = \left( a_{n}^{2s-1} - a_{n_{1}}^{2s-1} \right)^{2} + \left( a_{n}^{2s} - a_{n_{1}}^{2s} \right)^{2}, \\ \bar{a}_{2} = \left( a_{m}^{2s-1} - a_{m_{1}}^{2s-1} \right)^{2} + \left( a_{m}^{2s} - a_{m_{1}}^{2s} \right)^{2}$$

$$(28)$$

$$\bar{a}_3 = \left(a_n^{2s-1} - a_{n_1}^{2s-1}\right) \left(a_n^{2s} - a_{n_1}^{2s}\right) - \left(a_m^{2s-1} - a_{m_1}^{2s-1}\right) \left(a_m^{2s} - a_{m_1}^{2s}\right)$$
(29)

$$d_{n_1} = \left(\pi_1 P_d^{Q_{n_1}} (1 - P_d)^{K - Q_{n_1}} - \pi_0 P_f^{Q_{n_1}} (1 - P_f)^{K - Q_{n_1}}\right) \prod_{s=1}^S f\left(y_{2s-1}, y_{2s} | u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}\right)$$
(30)

where all entries of  $S_0$  and  $S_1$  are, respectively, negative and positive. We have

$$\begin{split} \mathcal{T}_{e_1|h} &< \frac{1}{|S_1|} \sum_{d_{n_1} \in S_1} P\left(|S_1|d_{n_1} > -d_n|F_n\right) \\ &= \frac{1}{|S_1|} \sum_{d_{n_1} \in S_1} P(\zeta(n,n_1) > -\ln(|S_1|G(n,n_1)) + I(n,n_1)) \\ \mathcal{T}_{e_2|h} &< \frac{1}{|S_0|} \sum_{d_{n_1} \in S_0} P\left(d_n < -|S_0|d_{n_1}|F_n\right) \\ &= \frac{1}{|S_0|} \sum_{d_{n_1} \in S_0} P(\zeta(n,n_1) > -\ln(|S_0|G(n,n_1)) + I(n,n_1)) \\ \mathcal{G}(n,n_1) &= \frac{\pi_1 P_d^{Q_{n_1}} (1 - P_d)^{K - Q_{n_1}} - \pi_0 P_f^{Q_{n_1}} (1 - P_f)^{K - Q_{n_1}}}{\pi_0 P_f^{Q_n} (1 - P_f)^{K - Q_n} - \pi_1 P_d^{Q_n} (1 - P_d)^{K - Q_n}} \\ I(n,n_1) &= \sum_{s=1}^{S} \mathcal{K}_s(n,n_1), \text{ where} \\ \mathcal{K}_s(n,n_1) &= \frac{1}{\sigma_v^2} \left( \left| h_{2s-1} \left( a_{n_1}^{2s-1} - a_n^{2s-1} \right) \right|^2 + \left| h_{2s} \left( a_{n_1}^{2s} - a_n^{2s} \right) \right|^2 \right) \\ \zeta(n,n_1) &= \sum_{s=1}^{S} \theta_s(n,n_1), \text{ where} \\ \theta_s(n,n_1) &= \frac{2}{\sigma_v^2} \Re \epsilon \left\{ v_{2s-1} h_{2s-1} \left( a_{n_1}^{2s-1} - a_n^{2s-1} \right) + v_{2s} h_{2s} \left( a_{n_1}^{2s} - a_n^{2s} \right) \right\}. \end{split}$$

Note that  $\zeta(n, n_1)$  is a zero mean Gaussian RV with the variance  $2I(n, n_1)$ . Using similar techniques in Section VI-A we can establish (32)–(35), shown at the bottom of the page.

This completes our derivations for the upper bounds on  $\overline{T}_{e_1}, \overline{T}_{e_2}$  and thus  $\overline{P}_{e_1}, \overline{P}_{e_2}$ . Our simulations indicate that the bound is minimized for  $t \approx 0.3$ .

# C. Cooperative Fusion Architecture With Signal Fusion at Sensors

Leveraging on the derivations in Section VI-A and B, we provide upper bounds on  $\overline{\mathcal{T}}_{e_1}, \overline{\mathcal{T}}_{e_2}$ . In fact, the absence of internode communication renders the notations in Section VI-A simple, as the indexes  $m, m_1$  and the decisions  $\hat{u}_{2s-1}, \hat{u}_{2s}$ are dropped,  $z_{2s-1}, z_{2s}$  are substituted with  $y_{2s-1}, y_{2s}$ , and the noise terms  $\delta_s^1, \delta_s^2$  are substituted with  $v_{2s-1}, v_{2s}$ . Thus, the derivations of the upper bounds become rather easy. In particular, instead of  $d_{n_1,m_1}$  in (21) or  $d_{n_1}$  in (30), we define  $d_{n_1}$  as (36), shown at the bottom of the page, where  $P(\tilde{u}_{2s-1} =$  $a_{n_1}^{2s-1}, \tilde{u}_{2s} = a_{n_1}^{2s} | \mathcal{H}_\ell)$  in (36) is determined in Section IV-B. In fact, this probability depends on sensing channels through threshold  $\tau$  and sensing noise variance  $\sigma_w^2$  as well as the average received SNR  $\bar{\gamma}_{hs}$  corresponding to internode communication. Furthermore, the relationship between  $P_{e_1|h}$ ,  $\mathcal{T}_{e_1|h}$  in (23) and  $P_{e_2|h}, \mathcal{T}_{e_2|h}$  in (24) can be revised as (31), in which  $F_n = \{\tilde{u}_{2s-1} = a_n^{2s-1}, \tilde{u}_{2s} = a_n^{2s}$  for  $s = 1, \ldots, S\}$ . We also redefine  $S_0 = \{d_{n_1} \text{ where } d_{n_1} < 0\}$  and  $S_1 = \{d_{n_1} \text{ where } d_{n_1} \ge 0\}$ , where all entries of  $S_0$  and  $S_1$  are, respectively, negative and positive. We have

$$\begin{split} \mathcal{T}_{e_1|h} &< \frac{1}{|S_1|} \sum_{d_{n_1} \in S_1} P\left(|S_1|d_{n_1} > -d_n|F_n\right) \\ &= \frac{1}{|S_1|} \sum_{d_{n_1} \in S_1} P(\zeta(n,n_1) > -\ln(|S_1|G(n,n_1)) + I(n,n_1)) \\ \mathcal{T}_{e_2|h} &< \frac{1}{|S_0|} \sum_{d_{n_1} \in S_0} P\left(d_n < -|S_0|d_{n_1}|F_n\right) \\ &= \frac{1}{|S_0|} \sum_{d_{n_1} \in S_0} P(\zeta(n,n_1) > -\ln(|S_0|G(n,n_1)) + I(n,n_1)) \end{split}$$

$$\bar{\mathcal{T}}_{e_1} < \frac{\mathbf{1}_{\{Q_n < M\}}}{2\sqrt{|S_1|}} \sum_{d_{n_1} \in S_1} \left[ \sqrt{G(n,n_1)} \prod_{s=1}^S \mathcal{D}_1(n,n_1) \right] + \mathbf{1}_{\{Q_n \ge M\}}$$
(32)

$$\bar{\mathcal{T}}_{e_2} < \frac{\mathbf{1}_{\{Q_n > M\}}}{|S_0|} \sum_{d_{n_1} \in S_0} \left[ \min_t \left( |S_0| G(n, n_1) \right)^t \prod_{s=1}^S \mathcal{D}_2(n, n_1) \right] + \mathbf{1}_{\{Q_n \le M\}}$$
(33)

$$\mathcal{D}_1(n,n_1) = \left( \left( 1 + \frac{\bar{\gamma}_h \left| a_n^{2s-1} - a_{n_1}^{2s-1} \right|}{2} \right) \left( 1 + \frac{\bar{\gamma}_h \left| a_n^{2s} - a_{n_1}^{2s} \right|}{2} \right) \right)^{-1}$$
(34)

$$\mathcal{D}_2(n,n_1) = \left( \left( 1 + \frac{4(t^2 - t)\bar{\gamma}_h \left| a_n^{2s-1} - a_{n_1}^{2s-1} \right|}{2} \right) \left( 1 + \frac{4(t^2 - t)\bar{\gamma}_h \left| a_n^{2s} - a_{n_1}^{2s} \right|}{2} \right) \right)^{-1}$$
(35)

$$d_{n_{1}} = \left(\pi_{1} \prod_{s=1}^{S} P\left(\tilde{u}_{2s-1} = a_{n_{1}}^{2s-1}, \tilde{u}_{2s} = a_{n_{1}}^{2s} | \mathcal{H}_{1}\right) - \pi_{0} \prod_{s=1}^{S} P\left(\tilde{u}_{2s-1} = a_{n_{1}}^{2s-1}, \tilde{u}_{2s} = a_{n_{1}}^{2s} | \mathcal{H}_{0}\right)\right) \times \prod_{s=1}^{S} f\left(y_{2s-1}, y_{2s} | \tilde{u}_{2s-1} = a_{n_{1}}^{2s-1}, \tilde{u}_{2s} = a_{n_{1}}^{2s}\right) \quad (36)$$

in which 
$$G(n, n_1) = (g(n_1)/(-g(n)))$$
  

$$g(n) = \pi_1 \prod_{s=1}^{S} P\left(\tilde{u}_{2s-1} = a_n^{2s-1}, \tilde{u}_{2s} = a_n^{2s} | \mathcal{H}_1\right)$$

$$-\pi_0 \prod_{s=1}^{S} P\left(\tilde{u}_{2s-1} = a_n^{2s-1}, \tilde{u}_{2s} = a_n^{2s} | \mathcal{H}_0\right)$$

$$I(n, n_1) = \sum_{s=1}^{S} \mathcal{K}_s(n, n_1), \text{ where}$$

$$\mathcal{K}_s(n, n_1) = \frac{\alpha}{\sigma_v^2} \left( |h_{2s-1}(a_{n_1}^{2s-1} - a_n^{2s-1})|^2 + |h_{2s}(a_{n_1}^{2s} - a_n^{2s})|^2 \right)$$

$$\zeta(n, n_1) = \sum_{s=1}^{S} \theta_s(n, n_1), \text{ where}$$

$$\theta_s(n, n_1) = \frac{2\sqrt{\alpha}}{\sigma_v^2} \Re \epsilon \left\{ v_{2s-1}h_{2s-1}\left(a_{n_1}^{2s-1} - a_n^{2s-1}\right) \right\}$$

 $+v_{2s}h_{2s}\left(a_{n_{1}}^{2s}-a_{n}^{2s}\right)\}.$ 

Note that  $\zeta(n, n_1)$  is a zero-mean Gaussian RV with variance  $2I(n, n_1)$ . Using similar techniques in Section VI-A and B, we can establish (37)–(40), shown at the bottom of the page.

This completes our derivations for the upper bounds on  $\overline{T}_{e_1}, \overline{T}_{e_2}$  and, thus,  $\overline{P}_{e_1}, \overline{P}_{e_2}$ . Our numerical results show that the bound is minimized for  $t \approx 0.3$ . Furthermore, the upper bounds depend on the power allocation parameter  $\alpha$ . In Section VII, we investigate the optimal  $\alpha$  that minimizes these bounds.

# D. Parallel Fusion Architecture With Local Threshold Changing at Sensors

Leveraging on the derivations in Section VI-A and B, we provide upper bounds on  $\bar{\mathcal{T}}_{e_1}, \bar{\mathcal{T}}_{e_2}$ . To capture all different values that  $u_{2s-1}, u_{2s}, \bar{u}_{2s-1}, \bar{u}_{2s}$  for  $s = 1, \ldots, S$  can take, we consider a 2K-length sequence  $(a_{n_1}^1, a_{n_1}^2, a_{m_1}^1, a_{m_2}^2, \ldots, a_{n_1}^{2S-1}, a_{n_1}^{2S}, a_{m_1}^{2S-1}, a_{m_1}^{2S})$ , where  $a_{n_1}^{2s-1}, a_{n_1}^{2s} \in \{1, -1\}$  and  $a_{m_1}^{2s-1}, a_{m_1}^{2s} \in \{1, -1\}$ , respectively, are the values assumed by  $u_{2s-1}, u_{2s}$  and  $\bar{u}_{2s-1}, \bar{u}_{2s}$ . Let  $Q_{n_1,m_1}^1, Q_{n_1,m_1}^2, Q_{n_1,m_1}^3, Q_{n_1,m_1}^4$ ,  $Q_{n_1,m_1}^4$ , respectively, denote the number of cases in the given sequence that  $a_{n_1}^{s'} = a_{m_1}^{s'} = 1$ ,  $a_{m_1}^{s'} = -a_{m_1}^{s'} = -1$ , and  $a_{n_1}^{s'} = a_{m_1}^{s'} = -1$  for s' = 2s, 2s - 1,

s = 1, ..., S. Instead of  $F_{n_1,m_1}$  in (17) and  $d_{n_1,m_1}$  in (18), we redefine them as follows:

$$F_{n_1,m_1} = \left\{ u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}, \bar{u}_{2s-1} = a_{m_1}^{2s-1}, \\ \bar{u}_{2s} = a_{m_1}^{2s} \text{ for } s = 1, \dots, S \right\}$$

$$d_{n_1,m_1} = \left( \pi_1 \prod_{j=1}^4 P_{d_j}^{Q_{n_1,m_1}^j} - \pi_0 \prod_{j=1}^4 P_{f_j}^{Q_{n_1,m_1}^j} \right) \\ \times \prod_{s=1}^S f\left( z_{2s-1}, z_{2s} | u_{2s-1} = a_{n_1}^{2s-1}, u_{2s} = a_{n_1}^{2s}, \\ \bar{u}_{2s-1} = a_{m_1}^{2s-1}, \bar{u}_{2s} = a_{m_1}^{2s} \right)$$
(41)

where 
$$\begin{split} P(\bar{u}_i = 1 | u_i = -u_j = -1, \mathcal{H}_\ell) &= P(x_i > \tau_1 | \mathcal{H}_\ell), \\ P(\bar{u}_i = -1 | u_i = -u_j = 1, \mathcal{H}_\ell) &= P(\tau < x_i < \tau_1 | \mathcal{H}_\ell), \\ P(\bar{u}_i = 1 | u_i = -u_j = -1, \mathcal{H}_\ell) &= P(\tau_2 < x_i < \tau | \mathcal{H}_\ell), \text{ and } \\ P(\bar{u}_i = -1 | u_i = -u_j = -1, \mathcal{H}_\ell) &= P(x_i < \tau_2 | \mathcal{H}_\ell), \text{ respectively, are equal to } P_{d_1}, P_{d_2}, P_{d_3}, P_{d_4} \text{ under } \mathcal{H}_1 \text{ and are equal to } P_{f_1}, P_{f_2}, P_{f_3}, P_{f_4} \text{ under } \mathcal{H}_0. \text{ Since sensing noise terms are identically distributed and uncorrelated, we find } \\ P(F_{n_1,m_1} | \mathcal{H}_1) &= \prod_{j=1}^4 P_{d_j}^{Q_{n_1,m_1}^j} \text{ and } P(F_{n_1,m_1} | \mathcal{H}_0) = \prod_{j=1}^4 P_{f_j}^{Q_{n_1,m_1}^j}. \text{ Note that the relationship between } P_{e_1|h}, \mathcal{T}_{e_1|h} \\ \text{ in (23) and } P_{e_2|h}, \mathcal{T}_{e_2|h} \text{ in (24) holds true. Using similar techniques in Section VI-A and B, we can establish (42)-(46), \\ \text{shown at the bottom of the next page.} \end{split}$$

This completes our derivations for the upper bounds on  $\overline{T}_{e_1}, \overline{T}_{e_2}$  and, thus,  $\overline{P}_{e_1}, \overline{P}_{e_2}$ . The upper bound on  $\overline{T}_{e_2}$  depends on t. Our simulations indicate that the bound is minimized for  $t \approx 0.3$ .

# *E.* Comparison of Different Schemes in Asymptotic Regime for Large S

For all four schemes discussed in Section VI-A–D, we have established the following, for large S (detailed derivations are available in [28, App. B]):

$$\begin{split} \bar{P}_{e} &= \bar{P}_{e_{1}} + \bar{P}_{e_{2}} < \kappa_{l_{11}} e^{S\left(\mu_{l_{11}} + \frac{1}{2}\sigma_{l_{11}}^{2}\right)} + \frac{1}{2} e^{-S\frac{\mu_{l_{12}}^{2}}{2\sigma_{l_{12}}^{2}}} \\ &+ \kappa_{l_{21}} e^{S\left(\mu_{l_{21}} + \frac{1}{2}\sigma_{l_{21}}^{2}\right)} + \frac{1}{2} e^{-S\frac{\mu_{l_{22}}^{2}}{2\sigma_{l_{22}}^{2}}} \\ \text{where } \mu_{l_{11}} + \frac{1}{2} \sigma_{l_{11}}^{2}, \mu_{l_{21}} + \frac{1}{2} \sigma_{l_{21}}^{2} < 0. \end{split}$$

$$\bar{\mathcal{T}}_{e_1} < \frac{\mathbf{1}_{\{d_n \in S_0\}}}{2\sqrt{|S_1|}} \sum_{d_{n_1} \in S_1} \left[ \sqrt{G(n,n_1)} \prod_{s=1}^S \mathcal{D}_1(n,n_1) \right] + \mathbf{1}_{\{d_n \in S_1\}}$$
(37)

$$\bar{\mathcal{T}}_{e_2} < \frac{\mathbf{1}_{\{d_n \in S_1\}}}{|S_0|} \sum_{d_{n_1} \in S_1} \left[ \min_t \left( |S_0| G(n, n_1) \right)^t \prod_{s=1}^S \mathcal{D}_2(n, n_1) \right] + \mathbf{1}_{\{d_n \in S_0\}}$$
(38)

$$\mathcal{D}_{1}(n,n_{1}) = \left( \left( 1 + \frac{\alpha \bar{\gamma}_{h} \left( a_{n}^{2s-1} - a_{n_{1}}^{2s-1} \right)^{2}}{4} \right) \left( 1 + \frac{\alpha \bar{\gamma}_{h} \left( a_{n}^{2s} - a_{n_{1}}^{2s} \right)^{2}}{4} \right) \right)^{-1}$$
(39)

$$\mathcal{D}_2(n,n_1) = \left( \left( 1 + \alpha (t^2 - t)\bar{\gamma}_h \left( a_n^{2s-1} - a_{n_1}^{2s-1} \right)^2 \right) \left( 1 + \alpha (t^2 - t)\bar{\gamma}_h \left( a_n^{2s} - a_{n_1}^{2s} \right)^2 \right) \right)^{-1}$$
(40)

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Moreover,  $\mu_{l_{11}}, \mu_{l_{12}}, \mu_{l_{21}}, \mu_{l_{22}}$  and  $\sigma_{l_{11}}^2, \sigma_{l_{12}}^2, \sigma_{l_{21}}^2, \sigma_{l_{22}}^2$  and  $\kappa_{l_{11}}, \kappa_{l_{21}}$  differ for different schemes and do not depend on S (they only depend on  $SNR_h$ ,  $SNR_c$  defined in Section VII and  $\pi_0$ ). For each scheme, we examine these four exponentials and keep the dominant exponential. Let  $\kappa_{l_{-}}e^{-S\gamma_{x}}$  for x = a, b, c, d be the dominant exponent, respectively, for schemes discussed in Section VI-A-D, where  $\gamma_x = \min\{-(\mu_{l_{11}} + (1/2)\sigma_{l_{11}}^2), (\mu_{l_{12}}^2/2\sigma_{l_{12}}), -(\mu_{l_{21}} +$  $(1/2)\sigma_{l_{21}}^2), (\mu_{l_{22}}^2/2\sigma_{l_{22}}^2)\}$ , and  $\kappa_{l_x}$  be its corresponding multiplicative scalar. When comparing the error exponents of any pair of these four schemes, for instance, schemes in Section VI-A and B, we have  $\lim_{S\to\infty}((\ln(\kappa_{l_a}e^{-S\gamma_a})/S) (\ln(\kappa_{l_b}e^{-S\gamma_b})/S)) = \gamma_b - \gamma_a$ , implying that such a difference depends on  $SNR_h, SNR_c, \pi_0$  only and does not change with S. This analysis suggests that our numerical findings in Section VII on performance comparison between different schemes should not vary much for large S = (K/2). In fact, our simulation results show that performance comparison between different schemes, given  $SNR_h$ ,  $SNR_c$ ,  $\pi_0$ , remains the same for K = 10, 14, 20 (due to space limitations, we only include the results for K = 20 in Table IV).

### VII. NUMERICAL RESULTS

Here, we evaluate and compare the performance of the proposed schemes in Sections III–V, against the conventional scheme in Section II-B. For the sake of presentation, we refer to the schemes in Sections II-B, III–V, respectively, as "parallel," "STC6@sensors," and "fusion6@sensors," "threshold changing6@sensors." We consider K=10 sensors (S=5 groups of paired sensors). We assume that the sensing noise terms

 $w_k$  are identically distributed, i.e.,  $\sigma_{w_k}^2 = \sigma_w^2$ , and  $\rho_{ij} = \rho$ characterizes the correlation. We define  $SNR_c = -20 \log_{10} \sigma_w$ as the SNR corresponding to sensing channels. We let the distances between the sensors and the FC d = 10 m, the distances between the cooperative partners within each group  $d_0 = 2$  m, the variance of receiver noise terms  $\sigma_v^2 = \sigma_\eta^2 = -50$  dBm, the path-loss exponent  $\varepsilon = 2$ , and the antenna gain  $\mathcal{G} = -30$  dB. To make a fair comparison among different schemes, we enforce the sensors in all schemes to transmit the same power  $\mathcal{P}$ . In "STC6@sensors" and "fusion6@sensors," a sensor spends  $(1 - \alpha)\mathcal{P}$  and  $\alpha\mathcal{P}$ , respectively, for communicating with its cooperative partner and with the FC, where  $\alpha$  is different in these two schemes. We define  $SNR_h = 10 \log_{10} \bar{\gamma}_h$ , in which  $\bar{\gamma}_h =$  $(\sigma_h^2/\sigma_v^2) = (P\mathcal{G}/d^{\varepsilon}\sigma_v^2)$ . Our goal is to investigate the average error  $\bar{P}_e$  of "STC6@sensors," "fusion6@sensors," and "threshold changing6@sensors" against that of "parallel," as SNR<sub>h</sub> and SNR<sub>c</sub> vary and identify different regimes in which these schemes outperform "parallel." Note that, in "STC6@sensors" and "fusion6@sensors," given  $SNR_h$  and  $SNR_c$ , average error  $\bar{P}_e$  depends on  $\alpha$ , i.e., one would expect that there is an optimal power allocation  $\alpha^*$  at which  $\bar{P}_e$  attains its minimum, given  $SNR_h$  and  $SNR_c$ . We start with investigating  $\alpha^*$ .

Optimal Power Allocation When the FC Employs LRT Rule: We start with "STC6@sensors." Fig. 1(a) plots  $\bar{P}_e$  versus  $\alpha$  for SNR<sub>h</sub> = 5 dB, SNR<sub>c</sub> = 2, 6, 10 dB and  $\pi_0 = 0.6$ , assuming  $\rho = 0$ . We observe that  $\alpha^* \approx 0.65$ , regardless of the variations in SNR<sub>c</sub> values. Fig. 1(b) plots  $\bar{P}_e$  versus  $\alpha$  for SNR<sub>c</sub> = 6 dB, SNR<sub>h</sub> = 5, 10, 15 dB and  $\pi_0 = 0.6$ , assuming  $\rho = 0$ . We observe that  $\alpha^*$  increases as SNR<sub>h</sub> (or, equivalently,  $\bar{\gamma}_h$ ) increases; in particular, we obtain  $\alpha^* \approx 0.6, 0.7, 0.8$ , respectively, for SNR<sub>h</sub> = 5, 10, 15 dB. These observations can be explained

$$\bar{\mathcal{T}}_{e_1} < \frac{\mathbf{1}_{\{d_{n,m}\in S_0\}}}{2\sqrt{|S_1|}} \sum_{d'_{n_1,m_1}\in S_1} \left[ \sqrt{G(n,m,n_1,m_1)} \prod_{s=1}^S \mathcal{D}_1(n,m,n_1,m_1) \right] + \mathbf{1}_{\{d_{n,m}\in S_1\}}$$
(42)

$$\bar{\mathcal{T}}_{e_2} < \frac{\mathbf{1}_{\{d_{n,m}\in S_1\}}}{|S_0|} \sum_{d_{n,m}\in S_0} \left[ \min_t (|S_0|G(n,m,n_1,m_1))^t \prod_{s=1}^S \mathcal{D}_2(n,m,n_1,m_1) \right] + \mathbf{1}_{\{d_{n,m}\in S_0\}}$$
(43)

$$\text{in which } S_{0} = \left\{ d_{n_{1},m_{1}} \text{ where } \frac{\prod_{j=1}^{4} P_{d_{j}}^{Q_{n_{1},m_{1}}}}{\prod_{j=1}^{4} P_{f_{j}}^{Q_{n_{1},m_{1}}}} < \frac{\pi_{0}}{\pi_{1}} \right\}, S_{1} = \left\{ d_{n_{1},m_{1}} \text{ where } \frac{\prod_{j=1}^{4} P_{d_{j}}^{Q_{n_{1},m_{1}}}}{\prod_{j=1}^{4} P_{d_{j}}^{Q_{n_{1},m_{1}}}} > \frac{\pi_{0}}{\pi_{1}} \right\}$$

$$(n,m,n_{1},m_{1}) = \frac{\left( \pi_{1} \prod_{j=1}^{4} P_{d_{j}}^{Q_{n_{1},m_{1}}} - \pi_{0} \prod_{j=1}^{4} P_{f_{j}}^{Q_{n_{1},m_{1}}} \right)}{\left( \pi_{1} \prod_{j=1}^{4} P_{d_{j}}^{Q_{n,m}} - \pi_{0} \prod_{j=1}^{4} P_{f_{j}}^{Q_{n,m}}} \right)}$$

$$(44)$$

$$\mathcal{D}_1(n, m, n_1, m_1) = \left( \left( 1 + \frac{\alpha \bar{\gamma}_h \bar{a}_1}{8} \right) \left( 1 + \frac{\alpha \bar{\gamma}_h \bar{a}_2}{8} \right) - \frac{\alpha^2 \bar{\gamma}_h^2 \bar{a}_3}{64} \right)^{-1}$$
(45)

$$\mathcal{D}_{2}(n,m,n_{1},m_{1}) = \left( \left( 1 + \frac{\alpha(t^{2}-t)\bar{\gamma}_{h}\bar{a}_{1}}{2} \right) \left( 1 + \frac{\alpha(t^{2}-t)\bar{\gamma}_{h}\bar{a}_{2}}{2} \right) - \frac{\alpha^{2}(t^{2}-t)^{2}\bar{\gamma}_{h}^{2}\bar{a}_{3}}{16} \right)^{-1}, \bar{a}_{1} = \left( a_{n}^{2s-1} - a_{n_{1}}^{2s-1} \right)^{2} + \left( a_{m}^{2s} - a_{m_{1}}^{2s} \right)^{2} \\ \bar{a}_{2} = \left( a_{m}^{2s-1} - a_{m_{1}}^{2s-1} \right)^{2} + \left( a_{n}^{2s} - a_{n_{1}}^{2s} \right)^{2} \\ \bar{a}_{3} = \left( a_{n}^{2s-1} - a_{n_{1}}^{2s-1} \right) \left( a_{m}^{2s} - a_{m_{1}}^{2s} \right) - \left( a_{m}^{2s-1} - a_{m_{1}}^{2s-1} \right) \left( a_{n}^{2s} - a_{n_{1}}^{2s} \right)$$

$$(46)$$



Fig. 1. "STC@sensors": LRT rule. (a)  $SNR_h = 5 dB$  with  $SNR_c = 2, 6, and 10 dB$ . (b)  $SNR_c = 6 dB$  with  $SNR_h = 5, 10, and 15 dB$ .



Fig. 2. "fusion@sensors": LRT rule. (a)  $SNR_h = 5 dB$  with  $SNR_c = 2$ , 6, and 10 dB. (b)  $SNR_c = 6 dB$  with  $SNR_h = 5$ , 10, and 15 dB.

considering our analytical results in (23)-(29) in Section VI-A. Recall that  $T_{n,m}, T_{n_1,m_1}$  in (27) capture the errors during internode communication and depend on the average received SNR  $\bar{\gamma}_{hs}$  corresponding to internode communication, which, for  $\sigma_{\eta}^2 = \sigma_v^2$ , reduces to  $\bar{\gamma}_{hs} = (d/d_0)^{\varepsilon} (1-\alpha) \bar{\gamma}_h$ . This implies that  $G(n, m, n_1, m_1)$  is decoupled into two fractions, where the first fraction depends on  $SNR_c$  (through the local performance indexes  $P_d, P_f$ ), and the second fraction depends on  $(d/d_0)^{\varepsilon}(1-\alpha)\bar{\gamma}_h$ . On the other hand, the inverses of  $\mathcal{D}_1(n, m, n_1, m_1), \mathcal{D}_2(n, m, n_1, m_1)$  depend on  $\alpha \bar{\gamma}_h$  only, capturing the errors of sensor-FC communication channels. Due to this decoupling of the effective factors in the terms of  $P_e$ , we expect that  $\alpha^*$  becomes insensitive to variations of SNR<sub>c</sub> (for fixed  $d, d_0$ , SNR<sub>h</sub>) and varies as SNR<sub>h</sub> changes (for fixed  $d, d_0, SNR_c$ ). For the scenario where the distance between cooperative partners is shorter than the distance between the nodes and the FC, we expect  $\alpha^* > 0.5$ , i.e., a sensor spends a higher (lower) percentage of its transmit power for communicating with the FC (its cooperative partner). These observations are in agreement with the fact that the local information exchange in "STC6@sensors" does not affect the local error probability  $p_{e_i} = P(u_i = -1|\mathcal{H}_1)\pi_1 + P(u_i = 1|\mathcal{H}_0)\pi_0$ at  $S_i$  (which depends on SNR<sub>c</sub> through  $P_d, P_f$ ); it rather provides a form of "decision diversity," to mitigate the fading effect during sensor-FC communication, i.e., it improves the global performance  $\bar{P}_e$  at the FC via reducing the errors during intersensor and sensor-FC communication. We continue with "fusion6@sensors." Fig. 2(a) plots  $\bar{P}_e$  versus  $\alpha$  for  $SNR_h = 5 dB$ ;  $SNR_c = 2$ , 6, and 10 dB; and  $\pi_0 = 0.6$ , assuming  $\rho = 0$ . We observe that  $\alpha^*$  increases as SNR<sub>c</sub> increases; in particular, we obtain  $\alpha^* \approx 0.6$ , 0.7, 0.8, respectively, for SNR<sub>c</sub> = 2, 6, 10 dB. Comparing this trend with that of "STC6@sensors" in Fig. 1(a), we notice that the schemes have different trends.

Fig. 2(b) plots  $\overline{P}_e$  versus  $\alpha$  for SNR<sub>c</sub> = 6 dB; SNR<sub>h</sub> = 5, 10, and 15 dB; and  $\pi_0 = 0.6$ , assuming  $\rho = 0$ . We observe that  $\alpha^*$  increases as SNR<sub>h</sub> (or, equivalently,  $\bar{\gamma}_h$ ) increases; in particular, we obtain  $\alpha^* \approx 0.7$ , 0.7, and 0.85, respectively, for  $SNR_h = 5$ , 10, and 15 dB. Comparing this trend with that of "STC6@sensors" in Fig. 1(b), we observe that the schemes have similar trends. These observations can be explained considering our analytical results in (37) and (38) in Section VI-C. Note that similar to "STC6@sensors," the inverses of  $\mathcal{D}_1(n, n_1), \mathcal{D}_2(n, n_1)$  depend on  $\alpha \overline{\gamma}_h$  only, capturing the errors of sensor-FC communication channels. However, the structure of  $G(n, n_1)$  is different from that of  $G(n, m, n_1, m_1)$ in "STC6@sensors." In particular, examining  $G(n, n_1)$  reveals that this term depends on  $P(\tilde{u}_{2s-1}, \tilde{u}_{2s} \text{ for all pairs} | \mathcal{H}_{\ell})$  given in (12), which, as we have previously mentioned in Section IV-B, depends on SNR<sub>c</sub> (through  $P_d, P_f$ ) as well as the average received SNR  $\bar{\gamma}_{hs}$  corresponding to internode communication. This implies that, different from  $G(n, m, n_1, m_1)$  in "STC6@sensors," the impacts of effective factors SNR<sub>c</sub> and  $(d/d_0)^{\varepsilon}(1-\alpha)\bar{\gamma}_h$  in  $G(n, n_1)$  cannot be decoupled, and hence,  $\alpha^*$  varies as SNR<sub>c</sub> changes (for fixed  $d, d_0, \text{SNR}_h$ ) or SNR<sub>h</sub> changes (for fixed  $d, d_0, SNR_c$ ). These observations are in agreement with the fact that the local information exchange in "fusion6@sensors," different from "STC6@sensors," affects the local error probability  $p_{e_i} = P(\tilde{u}_i = -1 | \mathcal{H}_1) \pi_1 +$  $P(\tilde{u}_i = 1 | \mathcal{H}_0) \pi_0$  at  $S_i$ . Therefore, it improves the global performance  $\bar{P}_e$  at the FC via improving the local error probability



Fig. 3. "STC@sensors": majority rule. (a)  $SNR_h = 5 dB$  with  $SNR_c = 2, 6$ , and 10 dB. (b)  $SNR_c = 6 dB$  with  $SNR_h = 5, 10$ , and 15 dB.



Fig. 4. "fusion@sensors": majority rule. (a)  $SNR_h = 5 dB$  with  $SNR_c = 2, 6, and 10 dB$ . (b)  $SNR_c = 6 dB$  with  $SNR_h = 5, 10, and 15 dB$ .

at the sensors. As SNR<sub>c</sub> decreases (for fixed  $d, d_0, \text{SNR}_h$ ), the reliability of the initial decision  $u_i$  at  $S_i$  (which is based on observation  $x_i$ ) reduces, and hence, the local information exchange is more needed to form the new decision  $\tilde{u}_i$  with higher reliability, where more local information exchange is translated into a higher (lower) percentage of transmit power for internode communication (sensor–FC communication) or, equivalently, smaller  $\alpha^*$ .

Optimal Power Allocation When the FC Employs Majority Rule: Similar observations are made when the FC employs the majority rule. Figs. 3(a) and 4(a), respectively, plot  $P_e$  versus  $\alpha$ , for "STC6@sensors" and "fusion6@sensors," when  $SNR_h =$ 5 dB,  $SNR_c = 2$ , 6, and 10 dB, and  $\pi_0 = 0.6$ . Figs. 3(b) and 4(b), respectively, plot  $\bar{P}_e$  versus  $\alpha$ , for "STC6@sensors" and "fusion6@sensors," when  $SNR_c = 6$  dB,  $SNR_h = 5$ , 10, and 15 dB, and  $\pi_0 = 0.6$ . Comparing Fig. 1(a) with Fig. 3(a), Fig. 2(a) with Fig. 4(a), Fig. 1(b) with Fig. 3(b), and Fig. 2(b) with Fig. 4(b), we can make similar observations regarding the variations of  $\alpha^*$  as SNR<sub>c</sub> or SNR<sub>h</sub> changes. For each scheme, when we compare the value of  $\alpha^*$  for LRT and majority rules, given  $d, d_0, \text{SNR}_c, \text{SNR}_h$ , we find that  $\alpha^*$  corresponding to the majority rule is larger than that of the LRT rule, i.e., a sensor spends a higher (lower) percentage of its transmit power for communicating with the FC (its cooperative partner). This is due to the fact that the majority rule demodulates first the sensor-FC channel outputs to find the channel inputs, rather than using the channel outputs directly for fusion, resembling the concept of "hard versus soft decoding" in [23]. To compensate for the information loss due to demodulation and its negative impact on error, each sensor is required to invest a



Fig. 5. Monte Carlo simulation versus analytical results.

higher percentage of its transmit power for communicating with the FC.

Performance Comparison of Different Schemes: To validate our performance analysis in Section VI, Fig. 5 shows  $\bar{P}_e$  of "parallel," "STC6@sensors," "fusion6@sensors," and "threshold changing6@sensors" versus SNR<sub>h</sub> for SNR<sub>c</sub> = 6 dB and  $\pi_0 = 0.6$ , to compare the analytical and Monte Carlo simulation results. We obtain  $\bar{P}_e$  of "STC6@sensors" and "fusion6@sensors," using  $\alpha^*$  corresponding to SNR<sub>h</sub> and SNR<sub>c</sub> values. The figure demonstrates a good agreement between theory and simulation. It also shows that, different from conventional communication systems,  $\bar{P}_e$  has an error floor at high SNR<sub>h</sub>. This behavior is due to the fact that  $\bar{P}_e$  in our distributed detection system is dependent on SNR<sub>h</sub> and SNR<sub>c</sub>.

	$SNR_c = 10dB$			S	$SNR_c = 6dI$	В	S	В		
SNR <sub>h</sub>	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
parallel	$1.9e{-4}$	3.8e - 5	$2.0e{-5}$	1.2e - 2	7.3e - 3	5.6e - 3	7.6e - 2	6.6e - 2	6.4e - 2	
STC	8.7e - 4	3.1e - 5	$1.4e{-5}$	1.8e - 2	6.9e - 3	5.4e - 3	8.9e-2	6.5e - 2	6.3e - 2	0.6
fusion	$1.9e{-4}$	1.7e - 5	$1.4e{-5}$	$1.2e{-2}$	4.8e - 3	3.9e - 3	7.4e-2	5.5e - 2	$5.2e{-2}$	$\pi_0 = 0.0$
threshold	$3.4e{-4}$	1.7e - 5	3.0e - 6	1.2e - 2	4.4e - 3	2.6e - 3	7.5e-2	4.7e - 2	$3.7e{-2}$	
parallel	3.0e - 4	4.3e - 5	$2.3e{-5}$	$1.3e{-2}$	7.4e - 3	6.7e - 3	8.1e-2	7.0e-2	6.6e - 2	
STC	8.9e - 4	2.9e - 5	$1.9e{-5}$	1.9e - 2	7.1e - 3	6.2e - 3	9.4e-2	6.9e - 2	6.5e - 2	0.7
fusion	$2.3e{-4}$	$2.0e{-5}$	7e-6	$1.3e{-2}$	5.9e - 3	5.4e - 3	7.9e-2	5.9e - 2	5.5e-2	$\pi_0 = 0.7$
threshold	$3.9e{-4}$	$1.5e{-5}$	0.0e - 6	$1.2e{-2}$	3.9e - 3	2.4e - 3	7.1e-2	4.6e - 2	$3.5e{-2}$	

TABLE I All Schemes, LRT Rule,  $\rho=0,\,K=10$ 

In fact, had all communication channels been error free,  $\bar{P}_e$  of "parallel" would be

$$\bar{P}_{e} = \pi_{1} P \left( \frac{P_{d}^{n} (1 - P_{d})^{2S-n}}{P_{f}^{n} (1 - P_{f})^{2S-n}} < \frac{\pi_{0}}{\pi_{1}} \right) + \pi_{0} P \left( \frac{P_{d}^{n} (1 - P_{d})^{2S-n}}{P_{f}^{n} (1 - P_{f})^{2S-n}} > \frac{\pi_{0}}{\pi_{1}} \right) = \pi_{1} \sum_{n=0}^{M-1} \frac{(2S)!}{(2S-n)!n!} P_{d}^{n} (1 - P_{d})^{2S-n} + \pi_{0} \sum_{n=M}^{2S} \frac{(2S)!}{(2S-n)!n!} P_{f}^{n} (1 - P_{f})^{2S-n}$$
(47)

where M satisfies  $(P_d^M(1-P_d)^{2S-M}/P_f^M(1-P_f)^{2S-M}) >$  $(\pi_0/\pi_1)$ . Equation (47) indicates that the error floor depends on  $SNR_c$  (through  $P_d, P_f$ ), and as  $SNR_c$  reduces, the error floor increases. Fig. 5 also shows that "parallel" and "STC6@sensors" have similar error floors, whereas "fusion6@sensors" have a lower error floor. These observations are in agreement with the fact that the local information exchange in "STC6@sensors" improves  $\bar{P}_e$  via providing "decision diversity," without changing the local error probability at  $S_i$  (which depends on the reliability of  $u_i$ ). On the other hand, the local information exchange in "fusion6@sensors" improves  $\bar{P}_e$  via improving the local error probability at  $S_i$  (which depends on the reliability of  $\tilde{u}_i$ ). For moderate/high  $SNR_h$ , where the errors during intersensor and sensor-FC communication are negligible,  $\bar{P}_e$  is governed by the local error probability at  $S_i$ . Since the reliability of local decisions in "parallel" and "STC6@sensors" is identical and the reliability of local decisions in "fusion6@sensors" exceeds that of "parallel" and "STC6@sensors," we expect that "parallel" and "STC6@sensors" have similar error floors, whereas "fusion6@sensors" reach a lower error floor, and Fig. 5 confirms these. Table I tabulates  $P_e$  of "parallel," "STC6@ sensors," "fusion6@sensors," and "threshold changing6@ sensors," as SNR<sub>h</sub> and SNR<sub>c</sub> vary, for  $\pi_0 = 0.6$  and  $\rho = 0$ , when the FC employs the LRT rule. To have a fair comparison among different schemes, we obtain  $\bar{P}_e$  of "STC6@sensors" and "fusion6@sensors," using  $\alpha^*$  corresponding to SNR<sub>h</sub> and SNR<sub>c</sub> values. Comparing "STC6@sensors" and "parallel," we note that for moderate  $SNR_h$  and moderate/high  $SNR_c$ , "STC@sensors" outperform "parallel," whereas the performance gain of "STC@sensors" decreases as SNR<sub>c</sub> reduces. On the other hand, for low  $SNR_h$ , "STC@sensors" perform

worse than "parallel," whereas for high SNR<sub>h</sub>, "parallel" and "STC@sensors" reach similar error floors. These observations agree with the fact that the local information exchange in "STC@sensors" improves  $\bar{P}_e$  via providing "decision diversity," without changing the local error probability at  $S_i$  (which depends on  $SNR_c$ ). For moderate/high  $SNR_h$ , the errors during intersensor and sensor-FC communication are small, and  $P_e$ is mainly determined by SNR<sub>c</sub>. Therefore, lowering SNR<sub>c</sub> increases  $\bar{P}_e$ . On the other hand, for low SNR<sub>h</sub>, the errors during intersensor communication negatively impact the diversity gain of "STC@sensors." Comparing "fusion@sensors" and "parallel," we note that for low SNR<sub>h</sub>, they have similar performance, whereas for moderate/high SNR<sub>h</sub>, "fusion@sensors" outperform "parallel" (regardless of  $SNR_c$ ). In particular, for high  $SNR_h$ , the error floor of "fusion@sensors" is smaller than that of "parallel." These observations agree with the fact that for moderate/high SNR<sub>h</sub>,  $\bar{P}_e$  is dominated by the local probability error at  $S_i$  and that the local probability error of "fusion@sensors" is smaller than that of "parallel." Comparing "threshold changing@sensors" and "parallel," we note that the local decisions in "threshold changing@sensors" have enhanced reliability, which is due to the fact that  $u_i, \bar{u}_i$  at  $S_i$  are obtained based on comparing the sensor's observation  $x_i$  with three thresholds (instead of one). For moderate/high  $SNR_h$ , "threshold changing@sensors" outperform "parallel" (regardless of  $SNR_c$ ). This observation can be explained as follows. In this SNR<sub>h</sub> regime,  $\bar{P}_e$  is dominated by the local probability error at  $S_i$ . Since the reliability of local decisions in "threshold changing@sensors" exceeds that of local decisions in "parallel," we expect that "threshold changing@sensors" outperform "parallel." Furthermore, "threshold changing@sensors" improve  $\bar{P}_e$  over "parallel," via providing "decision diversity" through Alamouti's STC. For low SNR<sub>h</sub>, "threshold changing@sensors" outperform "parallel" only for low SNR<sub>c</sub>. This is because in this regime, two factors contribute to  $P_e$ : unreliable local decisions and communication channel errors. Since "threshold changing@sensors" increase the reliability of local decisions, its performance exceeds that of "parallel." On the other hand, for low  $SNR_h$  and moderate/high  $SNR_c$ , where the communication channel errors are the major contributors to  $\bar{P}_e$ , Alamouti's STC introduces destructive signal interference and degrades the performance of "threshold changing@sensors" with respect to "parallel."

Table I also tabulates  $\bar{P}_e$  of "parallel," "STC@sensors," "fusion@sensors," and "threshold changing@sensors," as SNR<sub>h</sub> and SNR<sub>c</sub> vary, for  $\pi_0 = 0.7$  and  $\rho = 0$ , when the FC

	$SNR_c = 13dB$		$BdB$ $SNR_c = 10dB$			$SNR_c = 6dB$			S	В		
SNR <sub>h</sub>	5dB	10dB	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
parallel	2.4e-4	4e - 6	1.3e - 3	4.0e - 4	9.0e - 5	3.2e - 2	2.2e - 2	1.8e - 2	1.4e - 1	$1.3e{-1}$	$1.3e{-1}$	
STC	1.4e - 3	2e-6	4.2e - 3	2.3e - 4	$6.0e{-5}$	4.4e - 2	2.0e - 2	1.7e - 2	1.5e - 1	$1.3e{-1}$	$1.3e{-1}$	0.6
fusion	1.7e-4	3e-6	1.3e - 3	4.0e - 5	3e-5	2.7e-2	1.3e - 2	1.0e - 2	$1.2e{-1}$	9.7e - 2	9.0e-2	$\pi_0 = 0.0$
threshold			2.4e - 2	1.3e - 2	1.1e - 3	$1.2e{-1}$	$1.0e{-1}$	9.7e - 2	$2.3e{-1}$	$2.2e{-1}$	$2.1e{-1}$	
parallel			2.0e - 3	5.3e - 4	2.0e-4	6.1e-2	4.2e - 2	3.7e - 2	$2.0e{-1}$	$2.0e{-1}$	$2.0e{-1}$	
STC			4.9e - 3	4.7e - 4	1.6e - 5	6.8e - 2	3.9e - 2	3.6e - 2	1.9e - 1	$2.0e{-1}$	2.0e - 1	0.7
fusion			1.2e - 3	1.8e - 4	6e - 5	2.3e-2	1.4e-2	1.2e - 2	9.2e-2	1.0e - 1	9.8e - 2	$\pi_0 = 0.7$
threshold			2.2e-2	1.3e-2	1.0e - 3	$1.0e{-1}$	9.6e - 2	9.0e-2	$2.0e{-1}$	$1.9e{-1}$	$1.9e{-1}$	

TABLE II All Schemes, Majority Rule,  $\rho = 0, K = 10$ 

TABLE III All Schemes, LRT Rule,  $\pi_0 = 0.7, K = 10$ 

	S	$NR_c = 10d$	В		$SNR_c = 6dB$	3		$SNR_c = 2dE$	3	
SNR <sub>h</sub>	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
parallel	$7.4e{-4}$	5.3e - 4	3.8e - 4	2.8e - 2	2.4e - 2	2.2e-2	1.10e - 1	1.02e - 1	9.95e - 2	
STC	7.7e - 4	4.3e - 4	3.8e - 4	3.2e - 2	1.8e - 2	2.0e-2	1.18e - 1	1.02e - 1	9.89e - 2	a = 0.1
fusion	$6.9e{-4}$	3.0e-4	2.2e-4	2.5e - 2	1.7e - 2	1.6e - 2	1.05e - 1	9.5e-2	9.06e - 2	p = 0.1
threshold	$9.3e{-4}$	2.2e-4	1.5e-4	2.6e - 2	$1.5e{-2}$	$1.3e{-2}$	1.09e - 1	8.7e-2	7.81e-2	
parallel	2.7e - 3	1.9e - 3	1.7e - 3	4.0e - 2	3.6e - 2	3.5e - 2	1.38e - 1	1.32e - 1	1.31e - 1	
STC	3.4e - 3	1.8e - 3	1.6e - 3	4.4e - 2	3.5e - 2	3.5e - 3	1.43e - 1	1.32e - 1	1.32e - 1	
fusion	2.5e - 3	1.3e - 3	1.1e - 3	3.9e - 2	3.3e - 2	3.3e - 2	1.37e - 1	1.29e - 1	1.25e - 1	$\rho = 0.2$
threshold	3.1e - 2	1.4e - 3	1.0e - 3	$4.3e{-2}$	3.4e - 2	3.0e-2	1.42e - 1	1.29e - 1	1.24e - 1	
parallel	4.8e - 3	3.8e - 3	3.5e - 3	5.7e - 2	5.2e - 2	5.2e - 2	1.60e - 1	1.54e - 1	1.54e - 1	
STC	5.8e - 3	3.5e - 3	3.5e - 3	6.0e - 2	5.2e - 2	5.1e - 3	1.64e - 1	1.56e - 1	1.56e - 1	
fusion	4.8e - 3	3.4e - 3	3.3e - 3	5.6e - 2	5.0e - 2	4.9e - 2	1.60e - 1	1.52e - 1	1.52e - 1	p = 0.3
threshold	3.1e - 2	3.5e - 3	2.7e - 3	6.1e - 2	5.1e - 2	4.9e-2	1.62e - 1	1.54e - 1	1.50e - 1	
parallel	$1.4e{-2}$	1.4e-2	1.3e-2	9.19e - 2	8.95e - 2	8.99e - 2	1.80e - 1	1.77e - 1	1.77e - 1	
STC	$1.6e{-2}$	1.3e-2	1.3e-2	9.35e - 2	8.97e - 2	8.89e - 2	1.79e - 1	1.71e - 1	1.71e - 1	a = 0.5
fusion	1.4e - 3	1.2e-2	1.2e-2	9.12e - 2	8.79e - 2	8.55e-2	1.72e - 1	1.71e - 1	1.71e - 1	p = 0.5
threshold	$1.4e{-2}$	1.3e-2	1.0e-2	9.19e - 2	8.80e - 2	8.70e-2	1.80e - 1	1.78e - 1	1.77e - 1	
parallel	3.1e - 2	3.0e-2	2.9e-2	1.25e - 1	1.25e - 1	1.25e - 1	2.10e - 1	2.10e - 1	2.09e - 1	
STC	3.0e - 2	2.8e-2	2.8e-2	1.25e - 1	1.25e - 1	1.25e - 1	2.12e - 1	2.09e - 1	2.09e - 1	0 - 0 8
fusion	$2.9e{-2}$	2.9e-2	2.8e-2	1.25e - 1	1.25e - 1	1.25e - 1	2.09e - 1	2.09e - 1	2.09e - 1	p = 0.8
threshold	3.0e - 2	3.0e-2	3.0e-2	1.29e - 1	1.29e - 1	1.29e - 1	2.17e - 1	2.20e - 1	2.22e - 1	

employs the LRT rule. Similar observations can be made as we compare "STC@sensors" and "fusion@sensors" against "parallel" for different SNR<sub>h</sub> and SNR<sub>c</sub> regimes, regardless of  $\pi_0$ . However, when comparing "threshold changing@sensors" against "parallel" in low  $SNR_h$ , we note that the behavior changes as  $\pi_0$  increases. In particular, for low SNR<sub>h</sub>, "threshold changing@sensors" outperform "parallel" for low SNRc when  $\pi_0$  is smaller. As  $\pi_0$  increases, "threshold changing@sensors" outperform "parallel" for low/moderate SNR<sub>c</sub>, i.e., for low  $SNR_h$ , the range of  $SNR_c$  values over which "threshold changing@sensors" outperform "parallel" expands as  $\pi_0$  increases. Overall, Table I indicates that for low SNR<sub>h</sub> and moderate/high  $SNR_c$ , the proposed schemes do not have an advantage over "parallel." The exception is when  $\pi_0$  is large enough ( $\pi_0 \ge 0.7$ ), in which case "threshold changing@sensors" outperform "parallel." On the other hand, for low  $SNR_h$  and low  $SNR_c$ , "fusion@sensors" and "threshold changing@sensors" outperform "parallel." For moderate/high SNR<sub>h</sub>, regardless of SNR<sub>c</sub>, the schemes ranked from lowest to highest  $\bar{P}_e$  are "threshold changing@sensors," "fusion@sensors," "STC@sensors," and "parallel." Table II is similar to Table I, with the difference that the FC employs the majority rule. Comparing each of the schemes "STC@sensors," "fusion@sensors," and "threshold changing@sensors" against "parallel" for different SNR<sub>h</sub> and  $SNR_c$  regimes, we observe similar trends for the majority and LRT rules. However, when we compare the schemes to rank them based on their  $\bar{P}_e$ , we note the differences. In particular, for very high SNR<sub>c</sub>, "STC@sensors" outperform all the schemes, for high SNR<sub>c</sub>, none of the proposed schemes have an advantage over "parallel," and for moderate/low SNR<sub>c</sub>, "fusion@sensors" outperform all the schemes.

Impact of Correlation on Performance Comparison: Table III tabulates  $\bar{P}_e$  of "parallel," "STC@sensors," "fusion@sensors," and "threshold changing@sensors," as  $\text{SNR}_h$  and  $\text{SNR}_c$  vary, for  $\pi_0 = 0.7$  and  $\rho = 0.1, 0.2, 0.3,$ 0.5, 0.8, when the FC employs the LRT rule. We observe that as  $\rho$  increases, the performance gap between the proposed schemes and "parallel" reduces. This observation can be explained as follows. As we have previously mentioned, the performance advantage of "fusion@sensors" and "threshold changing@sensors" over "parallel," when the FC employs the LRT rules, is mainly due to the fact that the local information exchange in "fusion@sensors" or the three-threshold-based test at the sensors in "threshold changing@sensors" would enhance the reliability of the local decisions (compared with "parallel") when Gaussian sensing noise terms are uncorrelated. As these noise terms become correlated and  $\rho$  increases, the increase in the reliability of the local decisions in "fusion@sensors" and "threshold changing@sensors" diminishes, and thus, these two schemes start to lose their performance gain over "parallel." For  $\rho \leq 0.2$ , the observations made on the performance comparison among these schemes remain the

Π	5	$SNR_c = 100$	dB	S	$SNR_c = 6d$	В	$SNR_c = 2dB$			
SNR <sub>h</sub>	5dB	10dB	15dB	5dB	10dB	- 15dB	5dB	10dB		
parallel	3.8e-7	4.0e - 8	3.0e - 9	7.6e - 4	3.3e - 4	$2.3e{-4}$	2.3e - 2	1.6e - 2	1.5e-2	
STC	3.8e - 6	2.0e-8	2.0e-9	1.6e - 3	2.6e - 4	1.8e - 4	2.7e - 2	1.5e-2	1.3e - 2	
fusion	3.6e - 7	5.7e - 9	2.0e-9	6.0e - 4	1.7e-4	$9.5e{-5}$	2.1e - 2	1.3e-2	1.1e-2	
threshold	8.5e-7	6.0e - 9	$1.0e{-10}$	$7.3e{-4}$	1.1e-4	$3.8e{-5}$	2.1e - 2	9.1e - 3	6.0e - 3	

TABLE IV All Schemes, LRT Rule,  $\rho = 0, \pi_0 = 0.6, K = 20$ 

TABLE V All Schemes for a Group of Four Sensors, LRT Rule,  $\rho = 0, K = 4$ 

	$SNR_c = 10dB$			S	$SNR_c = 6dI$	3	S	3		
$SNR_h$	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
parallel	1.3e - 2	5.9e - 3	5.3e - 3	7.3e - 2	5.8e - 2	5.6e - 2	$1.8e{-1}$	$1.5e{-1}$	$1.5e{-1}$	
STC4	6.1e - 2	6.7e - 3	5.1e - 3	1.3e - 1	6.0e - 2	5.4e - 2	1.5e - 1	1.5e - 1	$1.5e{-1}$	
fusion4	1.3e - 2	2.8e - 3	1.6e - 3	7.8e - 2	4.5e - 2	3.7e - 2	$1.8e{-1}$	$1.4e{-1}$	$1.3e{-1}$	$\pi_0 = 0.0$
threshold4	$2.8e{-1}$	$2.4e{-1}$	$2.0e{-1}$	3.0e - 1	2.7e - 1	$1.9e{-1}$	$3.9e{-1}$	3.0e - 1	$2.2e{-1}$	
STC	2.1e - 2	5.8e - 3	5.1e - 3	8.8e - 2	5.7e - 2	5.4e - 2	$1.9e{-1}$	$1.5e{-1}$	$1.5e{-1}$	
fusion	1.3e - 2	3.8e - 3	3.6e - 3	7.4e-2	5.2e - 2	4.9e - 2	$1.8e{-1}$	1.5e - 1	$1.4e{-1}$	
threshold	1.5e - 2	2.8e - 3	1.5e - 3	7.8e - 2	4.4e - 2	3.5e - 2	$1.7e{-1}$	$1.3e{-1}$	$1.2e{-1}$	
parallel	1.0e - 2	5.5e - 3	4.8e - 3	6.7e - 2	5.3e - 2	5.0e - 2	$1.6e{-1}$	$1.5e{-1}$	$1.4e{-1}$	
STC4	5.7e - 2	8.3e - 3	3.5e - 3	$1.3e{-1}$	5.8e - 2	4.2e - 2	$2.0e{-1}$	$1.5e{-1}$	$1.4e{-1}$	$\pi_{2} = 0.7$
fusion4	$1.5e{-2}$	4.5e - 3	2.6e - 3	8.2e - 2	5.0e - 2	4.4e - 2	$1.7e{-1}$	$1.4e{-1}$	$1.3e{-1}$	$\pi_0 = 0.7$
threshold4	$2.3e{-1}$	$1.7e{-1}$	$2.0e{-1}$	$1.4e{-1}$	$2.0e{-1}$	$1.4e{-1}$	$2.9e{-1}$	$2.7e{-1}$	$1.2e{-1}$	
STC	1.8e - 2	5.0e - 3	4.1e - 3	7.9e - 2	5.0e - 2	$4.8e{-2}$	$1.7e{-1}$	$1.4e{-1}$	$1.4e{-1}$	
fusion	9.7e - 3	4.6e - 3	3.8e - 3	6.7e - 2	5.3e - 2	4.7e - 2	$1.6e{-1}$	$1.4e{-1}$	$1.4e{-1}$	
threshold	$1.2e{-2}$	3.1e - 3	2.2e - 3	7.0e-2	$4.5e{-2}$	$3.4e{-2}$	$1.6e{-1}$	$1.3e{-1}$	$1.2e{-1}$	

same as  $\rho = 0$ . When  $\rho$  varies between 0.2 and 0.3, "threshold changing@sensors" outperform others for high  $SNR_h$ , "fusion@sensors" outperform others for medium  $SNR_h$ , and "parallel" and "fusion@sensors" outperform others for low  $SNR_h$  (all regardless of  $SNR_c$ ). When  $\rho = 0.5$  for high  $SNR_h$ and high SNR<sub>c</sub>, "threshold changing@sensors" outperform others. For high SNR<sub>h</sub> and medium/low SNR<sub>c</sub> and for medium  $SNR_h$  (regardless of  $SNR_c$ ), "fusion@sensors" outperform others. For low  $SNR_h$  (regardless of  $SNR_c$ ), "parallel" and "fusion@sensors" outperform others. When  $\rho = 0.8$ , "threshold changing@sensors" have an inferior performance, regardless of  $SNR_h$  and  $SNR_c$ . Table III also shows that the performance degradation of "threshold changing@sensors" is pronounced, as  $\rho$  increases, compared with other schemes. Note that at  $\rho = 0$ , "threshold changing@sensors" have the lowest error floor, whereas at  $\rho = 1$ , all schemes have the same error floor. These imply that the rate of performance degradation of "threshold changing@sensors" must be higher than other schemes.

Impact of Increasing K: Table IV tabulates  $\overline{P}_e$  of "parallel," "STC@sensors," "fusion@sensors," and "threshold changing@sensors," as SNR<sub>h</sub> and SNR<sub>c</sub> vary, for  $\pi_0 = 0.6$ ,  $\rho = 0$ , and K = 20, when the FC employs the LRT rule. The observations made on the comparison between these schemes for K = 10 and 14 remain true.

Discussion on Increasing Number of Cooperative Partners in a Group: To investigate how the increasing number of partners impacts the performance, we consider a network of K = 4 sensors. Suppose sensors are positioned on the circumference of a circle on the x - y plane, whose center is located at the origin and its diameter is  $2\sqrt{2}$  m, and  $S_i$  is equally distant from  $S_j$  and  $S_k$  such that  $d_{ij} = d_{ik} = 2$  m, and  $d_{jk} = 2\sqrt{2}$  m. Moreover, the FC is located above the origin (above the x-y plane), such that all sensors are at equal distance of d = 10 m from the FC. Let "STC4@sensors," "fusion4@sensors," and "threshold changing4@sensors," respectively, refer to schemes 1)-3) with four partners in one group and "STC@sensors," "fusion@sensors," and "threshold changing@sensors," respectively, refer to schemes 1)-3) with two partners in one group (two groups in the network). Table V tabulates  $\bar{P}_e$  of all schemes as SNR<sub>h</sub> and SNR<sub>c</sub> vary, for  $\pi_0 = 0.6$ , 0.7 and  $\rho = 0$ , when the FC employs the LRT rule. Comparing all these schemes, we observe that for low SNR<sub>h</sub>, either "parallel" or "threshold changing@sensors" outperform others (depending on SNR<sub>c</sub>), whereas for moderate/ high SNR<sub>h</sub>, "threshold changing@sensors" outperform others, including "fusion4@sensors." These observations suggest that no performance gain is achieved as the number of cooperative partners increases beyond two. In the following, we provide our intuitive reasoning on how we expect schemes 1)–3) to perform, as the number of cooperative partners in a group increases, assuming K and sensor placements are fixed. Going from "STC@sensors" to "STC4@sensors," for fixed transmit power per sensor  $\mathcal{P}$ , we expect that the power consumption for intersensor communication  $(1 - \alpha)\mathcal{P}$ increases, as the average distances between sensors within a group increase. This leaves a sensor with smaller power, i.e., smaller  $\alpha \mathcal{P}$ , for its communication with the FC. For moderate/high SNR<sub>c</sub>, the relative performance of "STC@sensors" and "STC4@sensors" depends on the  $\mathcal{P}$  value. Our simulations show that for P > 32 mW,  $\alpha P$  is large enough that "STC4@sensors" provide a larger "decision diversity gain" than that of "STC@sensors" during sensor-FC communication, and thus, the former outperforms the latter. For low SNR<sub>c</sub>, however, these schemes have similar performance, since the performance in this regime is limited by the reliability of local decisions at sensors (which are the same for both schemes). Overall, these imply that for wireless sensor networks that

TABLE VI "Parallel," "Fusion@Sensors," and "Fusion6@Sensors," LRT Rule,  $\rho = 0, \pi_0 = 0.6, K = 6$ 

		S	$NR_c = 10d$	В	S	$SNR_c = 6dl$	В	$SNR_c = 2dB$			
	$SNR_h$	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
1	parallel	3.0e - 3	1.3e - 3	9.0e-4	3.7e-2	2.9e - 2	2.7e - 2	$1.3e{-1}$	$1.1e{-1}$	$1.1e{-1}$	
f	fusion	2.8e - 3	8.0e-4	3.0e - 4	3.5e - 2	2.1e - 2	1.7e - 2	$1.2e{-1}$	9.5e - 2	9.0e - 2	
f	fusion6	3.1e - 3	1.3e - 3	3.1e-4	3.8e - 2	2.9e - 2	2.6e - 2	$1.4e{-1}$	$1.1e{-1}$	$9.9e{-2}$	

TABLE VII All Schemes, LRT Rule,  $\rho = 0, \pi_0 = 0.6, K = 4$ 

	$SNR_c = 10dB$			5	$SNR_c = 6dl$	В	5	В		
SNR <sub>h</sub>	5dB	10dB	15dB	5dB	10dB	15dB	5dB	10dB	15dB	
parallel	1.3e - 2	5.9e - 3	5.3e - 3	7.3e - 2	5.8e - 2	5.6e - 2	1.8e - 1	1.5e - 1	1.5e - 1	
STC	2.1e-2	5.8e - 3	5.1e - 3	8.8e - 2	5.7e-2	5.4e - 2	1.9e - 1	1.5e - 1	1.5e - 1	homogeneus
fusion	1.3e - 2	3.8e - 3	3.6e - 3	7.4e-2	5.2e-2	4.9e - 2	1.8e - 1	1.5e - 1	1.4e - 1	nomogenous
threshold	1.5e - 2	2.8e - 3	1.5e - 3	7.8e - 2	4.4e-2	3.5e - 2	1.7e - 1	$1.3e{-1}$	$1.2e{-1}$	
$\mathcal{P}$	3.2mW	10mW	32mW	3.2mW	10mW	32mW	3.2mW	10mW	32mW	
parallel	1.1e-2	6.4e - 3	5.2e - 3	7.5e-2	6.0e-2	5.8e - 2	1.7e - 1	1.5e - 1	1.5e - 1	
STC	1.4e-2	6.9e - 3	5.1e - 3	8.4e - 2	5.8e - 2	5.6e - 2	1.8e - 1	1.5e - 1	1.5e - 1	inhoneoconous
fusion	1.1e - 2	4.0e - 3	3.6e - 3	7.3e-2	5.4e-2	4.9e - 2	1.7e - 1	1.5e - 1	1.4e - 1	mnomogenous
threshold	1.0e-2	3.6e - 3	1.4e-3	6.3e - 2	4.5e-2	3.3e - 2	1.6e - 1	$1.4e{-1}$	1.2e - 1	

typically operate within  $0.12 \le \mathcal{P} \le 36$  mW [27], increasing the number of cooperative partners in a group beyond two does not have much practical incentive. Similarly, going from "fusion@sensors" to "fusion4@sensors," we expect that  $(1 - \alpha)\mathcal{P}$  increases, whereas  $\alpha \mathcal{P}$  decreases. However, different from scheme 1), in scheme 2), local information exchange affects the reliability of local decisions. Hence, the increase in the reliability of the local decisions in "fusion4@sensors" due to the increase of  $(1 - \alpha)\mathcal{P}$  still compensates for a less reliable sensor-FC communication due to the decrease of  $\alpha \mathcal{P}$ , which leads to the observation that "fusion4@sensors" outperform "fusion@sensors" for  $\mathcal{P} \ge 10$  mW. Going to "fusion6@sensors" for K = 6, however,  $\alpha$  decreases further, such that even for a large  $\mathcal{P} \approx 40$  mW, the unreliability of the sensor-FC communication due to the decrease of  $\alpha P$  leads to the observation that "fusion6@sensors" perform worse than "fusion@sensors" (see Table VI). We expect similar performance degradation as we increase the number of partners in a group beyond six. Going from "threshold changing@sensors" to "threshold changing4@sensors," the former outperforms the latter for all  $SNR_c$  and  $SNR_h$ . This is due to the fact that as the number of cooperative partners in a group increases, the chance that the corresponding information matrix transmitted by this group to the FC deviates from the conventional orthogonal STC matrix increases, leading to destructive signal interference at the FC and diminishing the "decision diversity gain" of STC. We conjecture similar performance degradation as we increase the number of partners in a group beyond four.

Homogeneous Versus Inhomogeneous Sensor Placement: We consider a network of K = 4 sensors, consisting of two groups, where sensors are positioned on the circumference of a circle on the x-y plane, whose center is located at the origin and its diameter is 20 m. The distance between two sensors in each group is  $d_0 = 2$  m. For homogeneous placement, we assume that the FC is located at the origin, and for inhomogeneous placement, we move the FC toward one of the groups, such that the distance between the FC and the two groups are 4 and 16 m. Table VII tabulates  $\bar{P}_e$  of "parallel," "STC@sensors," "fusion@sensors," and "threshold changing@sensors," as SNR<sub>h</sub> and SNR<sub>c</sub> vary, for  $\pi_0 = 0.6$  and  $\rho = 0$ , when the FC employs the LRT rule, for homogeneous and inhomogeneous placements. This table shows that our findings on comparison between different schemes are exactly the same as those in Table I, which was another example of a homogenous placement. On the other hand, for inhomogeneous placement, "threshold changing@sensors" outperform other schemes, regardless of  $\mathcal{P}$  and SNR $_c$  values. Since "threshold changing@sensors" have the lowest error floor among all schemes, when one group of sensors becomes closer to the FC, the enhanced reliability of the information delivered to the FC by this group leads to the superior performance of this scheme. Comparing "STC@sensors" and "parallel," we note that the former performs worse than the latter in inhomogeneous placement. Since "STC@sensors" and "parallel" have similar error floor, placing one group of sensors closer to the FC does not change the reliability of the information provided by this group to the FC in either schemes. However, since the other group of sensors becomes farther away from the FC, the quality of the information delivered to the FC by this group decreases (due to destructive signal interference of STC), leading to the inferior performance of "STC@sensors."

### VIII. CONCLUSION

For the problem of binary distributed detection in a wireless sensor network, we have proposed novel cooperative and parallel fusion architectures, to combat the fading effects encountered in the conventional parallel fusion architecture. In particular, we have proposed 1) cooperative fusion architecture with Alamouti's STC scheme at sensors, 2) cooperative fusion architecture with signal fusion at sensors, and 3) parallel fusion architecture with local threshold changing at sensors. While there is limited local information exchange among the sensors (1-bit message) in schemes 1) and 2), there is no explicit information exchange in scheme 3). For these schemes, we derived the optimal LRT and the suboptimal majority fusion rules and analyzed their performance, in terms of communication and sensing SNRs. Our numerical results show that when the FC employs the LRT rule, unless for low communication SNR and moderate/high sensing SNR, performance improvement is feasible with the new cooperative and parallel fusion architectures, while scheme 3) outperforms the others. When the FC utilizes the majority rule, such improvement is possible, unless for high sensing SNR. In particular, for very high sensing SNR, scheme 1) outperforms the others, whereas for moderate/low sensing SNR, scheme 2) outperforms the others.

# APPENDIX

Upper Bounds on  $\mathcal{T}_{e_1|h}$  in (23) and its average  $\overline{\mathcal{T}}_{e_1}$ : For  $Q_n < M$ , we have

$$\mathcal{T}_{e_1|h} = P\left(\sum_{d_{n_1,m_1}\in S_1} d_{n_1,m_1} > -\sum_{d_{n_1,m_1}\in S_0} d_{n_1,m_1}|F_{n,m}\right)$$
$$< P\left(\sum_{d_{n_1,m_1}\in S_1} d_{n_1,m_1} > -d_{n,m}|F_{n,m}\right) \quad (48)$$

where the bound in (48) is obtained, noting that  $d_{n,m} \in S_0$ and that  $\sum_{d_{n_1,m_1} \in S_0} d_{n_1,m_1} < d_{n,m}$ . To further bound (48), we define interval x and function  $\varphi$ , as follows:

$$x = \mathbb{E}\left\{\sum_{d_{n_1,m_1}\in S_1} C_{n_1,m_1} d_{n_1,m_1} | F_{n,m}\right\}$$
$$\varphi(x) = P\left(\sum_{d_{n_1,m_1}\in S_1} C_{n_1,m_1} d_{n_1,m_1} > -d_{n,m} | F_{n,m}\right)$$
(49)

where constants  $C_{n_1,m_1}$  take values in the interval  $[0, |S_1|]$ . Our numerical results suggest that for small  $|S_1|$ ,  $\varphi$  is convex over x. Invoking the inequality  $\varphi(\sum_{i=1}^n x_i/n) \leq (\sum_{i=1}^n \varphi(x_i)/n)$ , where the points  $x_1, \ldots, x_n$  belong to x [26], and letting  $n = |S_1|$  and  $x_i = \mathbb{E}\{|S_1|d_{n_1,m_1}|F_{n,m}\}$  for  $i = 1, \ldots, |S_1|$ , we establish the following:

$$\varphi\left(\sum_{d_{n_1,m_1}\in S_1} \mathbb{E}\{d_{n_1,m_1}|F_{n,m}\}\right)$$
  
=  $\varphi\left(\frac{1}{|S_1|}\sum_{d_{n_1,m_1}\in S_1} \mathbb{E}\{|S_1|d_{n_1,m_1}|F_{n,m}\}\right)$   
 $\leq \frac{1}{|S_1|}\sum_{d_{n_1,m_1}\in S_1} \varphi\left(\mathbb{E}\{|S_1|d_{n_1,m_1}|F_{n,m}\}\right).$  (50)

The inequality in (50) implies that the upper bound on  $\mathcal{T}_{e_1|h}$  in (48) can be further bounded as

$$P\left(\sum_{d_{n_1,m_1}\in S_1} d_{n_1,m_1} > -d_{n,m}|F_{n,m}\right)$$
  
$$\leq \frac{1}{|S_1|} \sum_{d_{n_1,m_1}\in S_1} P\left(|S_1|d_{n_1,m_1} > -d_{n,m}|F_{n,m}\right). \quad (51)$$

The new bound on  $\mathcal{T}_{e_1|h}$  in (51) can be presented in closed form, considering the definitions of  $d_{n,m}$  and  $d_{n_1,m_1}$  in (21) and noting that, conditioned on  $u_{2s-1}, u_{2s}, \hat{u}_{2s-1}, \hat{u}_{2s}$ , the terms  $z_{2s-1}, z_{2s}$  are independent complex Gaussian RVs with the variance  $\sigma^2 = (|h_{2s-1}|^2 + |h_{2s}|^2)\sigma_v^2$  and the mean values  $\mu_{2s-1}^{n_1,m_1}, \mu_{2s}^{n_1,m_1}$  for  $d_{n_1,m_1}$  and  $\mu_{2s-1}^{n,m}, \mu_{2s}^{n,m}$  for  $d_{n,m}$ . Mapping the noise terms  $\delta_{ij}^1, \delta_{ij}^2$  in Section III-B into  $\delta_s^1, \delta_s^2$ , we find

$$P(|S_1|d_{n_1,m_1} > -d_{n,m}|F_{n,m}) = P(\zeta(n,m,n_1,m_1))$$
  
> -ln(|S\_1|G(n,m,n\_1,m\_1)) + I(n,m,n\_1,m\_1)) (52)

where  $G(n, m, n_1, m_1)$  is defined in (27), and  $I(n, m, n_1, m_1) = \sum_{s=1}^{S} \mathcal{K}_s(n, m, n_1, m_1)$  in which

$$\begin{aligned} \mathcal{K}_{s}(n,m,n_{1},m_{1}) \\ &= \frac{1}{\sigma^{2}} \left( \left| \mu_{2s-1}^{n_{1},m_{1}} - \mu_{2s-1}^{n,m} \right|^{2} + \left| \mu_{2s}^{n_{1},m_{1}} - \mu_{2s}^{n,m} \right|^{2} \right) \\ &= \frac{\alpha}{2\sigma_{v}^{2}} \left( \left| h_{2s-1}(a_{n_{1}}^{2s-1} - a_{n}^{2s-1}) + h_{2s}\left(a_{n_{1}}^{2s} - a_{n}^{2s}\right) \right|^{2} \\ &+ \left| h_{2s-1}\left(a_{m_{1}}^{2s-1} - a_{m}^{2s-1}\right) + h_{2s}\left(a_{m_{1}}^{2s} - a_{m}^{2s}\right) \right|^{2} \right) \end{aligned}$$

$$(53)$$

 $\zeta(n,m,n_1,m_1)$ 

$$= \sum_{s=1}^{S} \theta_{s}(n, m, n_{1}, m_{1}) \quad \text{where} \quad \theta_{s}(n, m, n_{1}, m_{1})$$
$$= \frac{2}{\sigma^{2}} \Re \mathfrak{e} \left\{ \delta_{s}^{1} \left( \mu_{2s-1}^{n_{1}, m_{1}} - \mu_{2s-1}^{n, m} \right)^{*} + \delta_{s}^{2} \left( \mu_{2s}^{n_{1}, m_{1}} - \mu_{2s}^{n, m} \right)^{*} \right\}.$$
(54)

Recall that  $\delta_s^1$  and  $\delta_s^2$  are i.i.d. zero-mean complex Gaussian RVs with variance  $\sigma^2$ . Hence,  $\zeta(n, m, n_1, m_1)$  is a zero-mean Gaussian RV with variance  $2I(n, m, n_1, m_1)$ . Thus, we can express (52) as

$$P(|S_1|d_{n_1,m_1} > -d_{n,m}|F_{n,m}) = Q\left(\frac{-\ln(|S_1|G(n,m,n_1,m_1)) + I(n,m,n_1,m_1)}{\sqrt{2I(n,m,n_1,m_1)}}\right).$$
(55)

Note that  $I(n,m,n_1,m_1)$  depends on the coefficients  $h_{2s-1}, h_{2s}$ , whereas  $G(n,m,n_1,m_1)$  is independent of these coefficients. In fact,  $G(n,m,n_1,m_1)$  depends on sensing channels through  $P_d, P_f$  and the average received SNR  $\bar{\gamma}_{hs}$  corresponding to internode communication through  $T_{n,m}, T_{n_1,m_1}$ . One can verify that when  $\pi_0 > \pi_1$ , we have  $-\ln(|S_1|G(n,m,n_1,m_1)) + I(n,m,n_1,m_1) > 0$ . Combining (48), (51), and (55), using the Chernoff bound of Q-function  $Q(x) < (1/2)e^{-(x^2/2)}$  for x > 0 and noting that  $0 < e^{(-(\ln(|S_1|G(n,m,n_1,m_1)))^2/4I(n,m,n_1,m_1))} < 1$  and, thus, can be dropped without decreasing the upper bound, we find

$$\mathcal{T}_{e_1|h} < \frac{\mathbf{1}_{\{Q_n < M\}}}{2\sqrt{|S_1|}} \sum_{d_{n_1,m_1} \in S_1} \sqrt{G(n,m,n_1,m_1)} \\ \times e^{-\frac{I(n,m,n_1,m_1)}{4}} + \mathbf{1}_{\{Q_n \ge M\}}.$$
 (56)

Finally, to find an upper bound on  $\overline{\mathcal{T}}_{e_1}$ , we need to take the average of  $e^{-(I(n,m,n_1,m_1)/4)}$  in (56) over  $h_{2s-1}, h_{2s}$  for  $s = 1, \ldots, S$ . Since  $h_{2s-1}, h_{2s} \sim \mathcal{CN}(0, \sigma_h^2)$  are i.i.d. across the pairs, we have

$$\mathbb{E}\left\{e^{-\frac{I(n,m,n_{1},m_{1})}{4}}\right\}$$

$$=\prod_{s=1}^{S}\int_{h_{2s-1}}\int_{h_{2s}}e^{-\frac{\mathcal{K}_{s}(n,m,n_{1},m_{1})}{4}}e^{-\frac{\left(|h_{2s-1}|^{2}+|h_{2s}|^{2}\right)}{\sigma_{h}^{2}}}d_{h_{2s-1}}d_{h_{2s}}.$$
(57)

After some calculations, (57) is reduced to  $\prod_{s=1}^{S} \mathcal{D}_1(n, m, n_1, m_1)$ , where  $\mathcal{D}_1(n, m, n_1, m_1)$  is given in (28). The upper bound on  $\overline{\mathcal{T}}_{e_1}$  is obtained by substituting  $e^{-(I(n,m,n_1,m_1)/4)}$  in (56) with  $\prod_{s=1}^{S} \mathcal{D}_1(n,m,n_1,m_1)$ . This completes our derivations for the upper bound on  $\overline{\mathcal{T}}_{e_1}$ .

Upper Bounds on  $\mathcal{T}_{e_2|h}$  in (24) and its average  $\overline{\mathcal{T}}_{e_2}$ : For  $Q_n > M$ , we have

$$\mathcal{T}_{e_2|h} < P(d_{n,m} < -\sum_{d_{n_1,m_1} \in S_0} d_{n_1,m_1} | F_{n,m}) < \frac{1}{|S_0|} \sum_{d_{n_1,m_1} \in S_0} P(d_{n,m} < -|S_0| d_{n_1,m_1} | F_{n,m})$$
(58)

noting that  $d_{n,m} \in S_1$  and  $\sum_{d_{n_1,m_1} \in S_1} d_{n_1,m_1} > d_{n,m}$ . The new bound on  $\mathcal{T}_{e_2|h}$  in (58) can be found in closed form, via examining the definitions of  $d_{n,m}$  and  $d_{n_1,m_1}$  in (21) and noting that, conditioned on  $u_{2s-1}, u_{2s}, \hat{u}_{2s-1}, \hat{u}_{2s}$ , the terms  $z_{2s-1}, z_{2s}$  are independent complex Gaussian RVs with variance  $\sigma^2 = (|h_{2s-1}|^2 + |h_{2s}|^2)\sigma_v^2$  and mean values  $\mu_{2s-1}^{n_1,m_1}, \mu_{2s}^{n_1,m_1}$  for  $d_{n_1,m_1}$  and  $\mu_{2s-1}^{n,m}, \mu_{2s}^{n,m}$  for  $d_{n,m}$ . Therefore

$$P(d_{n,m} < -|S_0|d_{n_1,m_1}|F_{n,m}) = P(\zeta(n,m,n_1,m_1))$$
  
> -\ln(|S\_0|G(n,m,n\_1,m\_1)) + I(n,m,n\_1,m\_1)). (59)

Comparing (52) and (59), it seems natural to write (59) in terms of the Q-function and apply the Chernoff bound to reach a bound. However, different from (52), when  $\pi_0 > \pi_1$ , we no longer have  $-\ln(|S_0|G(n,m,n_1,m_1)) + I(n,m,n_1,m_1) > 0$ . We use an alternative bound, which states  $P(\sum_{s=1}^{S} x_s < a) < \min_{t,t>0} e^{ta} \prod_{s=1}^{S} \mathbb{E}\{e^{-tx_s}\}$  when  $x_1, \ldots, x_S$  are independent RVs [26]. Combining (58) and (59), letting  $x_s = -\theta_s(n,m,n_1,m_1)$ ,  $a = \ln(|S_0|G(n,m,n_1,m_1)) - I(n,m,n_1,m_1)$  in (59), and using the alternative bound, we find

$$\mathcal{T}_{e_{2}|h} < \frac{\mathbf{1}_{\{Q_{n} > M\}}}{|S_{0}|} \sum_{d_{n_{1},m_{1}} \in S_{0}} \left[ \min_{t} e^{-tI(n,m,n_{1},m_{1})} \times (|S_{0}|G(n,m,n_{1},m_{1}))^{t} \prod_{s=1}^{S} \mathbb{E} \left\{ e^{t\theta_{s}(n,m,n_{1},m_{1})} \right\} \right] + \mathbf{1}_{\{Q_{n} \le M\}}$$

$$(60)$$

Noting that  $-\theta_s(n, m, n_1, m_1) \sim C\mathcal{N}(0, 2\mathcal{K}_s(n, m, n_1, m_1))$ and using the moment-generating function results, we find  $\mathbb{E}\left\{e^{t\theta_s(n,m,n_1,m_1)}\right\} = e^{t^2 \mathcal{K}_s(n,m,n_1,m_1)}$ . This implies that we can rewrite (60) as follows:

$$\mathcal{T}_{e_{2}|h} < \frac{\mathbf{1}_{\{Q_{n} > M\}}}{|S_{0}|} \sum_{d_{n_{1},m_{1}} \in S_{0}} \left[ \min_{t} (|S_{0}|G(n,m,n_{1},m_{1}))^{t} \\ \times \prod_{s=1}^{S} e^{(t^{2}-t)\mathcal{K}_{s}(n,m,n_{1},m_{1})} \right] + \mathbf{1}_{\{Q_{n} \le M\}}.$$
 (61)

To find a bound on  $\overline{\mathcal{T}}_{e_2}$ , we need to take the average of  $e^{(t^2-t)\mathcal{K}_s(n,m,n_1,m_1)}$  in (61) over  $h_{2s-1}, h_{2s}$ . One can verify that this term is equal to  $\mathcal{D}_2(n,m,n_1,m_1)$  in (29). The upper bound on  $\overline{\mathcal{T}}_{e_2}$  is obtained by substituting  $e^{(t^2-t)\mathcal{K}_s(n,m,n_1,m_1)}$  in (61) with (29). This completes our derivations for the bound on  $\overline{\mathcal{T}}_{e_2}$ .

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